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RECONSTRUCTING EQUIVALENCES

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Abstract: A graph is called an equivalence if each of its components of connectivity is a complete graph. We ask whether an equivalence is uniquely determined with its k -point subobjects. For each k we prove: 1/Every equivalence on less than $k \cdot \ln(k/2)$ -point set is uniquely determined with k -point subobjects; 2/It is not true that every equivalence on at least $(k+1) \cdot 2^{k-1}$ -point set is uniquely determined with its k -point subobjects.

0. Introduction

We denote $\langle V, W \rangle$ the ordered pair where the first member is V and the second one is W . $P_2(X)$ denotes the set of all 2-point subsets of the set X . An ordered pair $G = \langle X, R \rangle$ where $R \subset P_2(X)$ is called a graph and we denote $|G| = \text{card } X$ the number of points of X . The complete graph on X is the graph $\langle X, P_2(X) \rangle$ and we denote K_n the standard complete graph on n -point set. For the graph $G = \langle X, R \rangle$ and the set $Y \subset X$ we define the induced graph $G/Y = \langle Y, R \cap P_2(Y) \rangle$. In usual sense we work with concepts in graph theory, namely the connectivity of graphs, components of connectivity, isomorphism of graphs. The number of components of G is denoted $\text{cp } G$; isomorphic graphs are denoted $G \simeq H$ and nonisomorphic graphs $G \neq H$.

For every sequence of complete graphs K_{n_1}, \dots, K_{n_s} it is the standard sum $K = K_{n_1} + \dots + K_{n_s}$ with components of connectivity C_1, \dots, C_s satisfying $K/C_i \simeq K_{n_i}$; if $n_1 = \dots = n_s = n$ we write simply $K = s \cdot K_n$.

Definition 0.1. A graph E is called an equivalence if E is isomorphic to a sum of complete graphs.

Definition 0.2. The frequency of the graph H in the graph G

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is the number $\text{frq}(H, G) = \text{card} \{Y; G/Y \simeq H\}$. For an integer k the notation $G_1 \overset{k}{\sim} G_2 / G_1 \overset{\leq k}{\sim} G_2$, respectively/ means that for every graph H such that $|H| = k / |H| \leq k$, respectively/ the equality $\text{frq}(H, G_1) = \text{frq}(H, G_2)$ holds.

Remark 0.3. An induced graph of an equivalence is an equivalence. Thus, if E is an equivalence then $\text{frq}(H, E) > 0$ if and only if H is an equivalence.

We have showed in [6] the following theorem.

Theorem 0.4. Let k be an integer, G_1, G_2 be graphs. Following three properties are equivalent /i/ $G_1 \overset{k}{\sim} G_2$, /ii/ $G_1 \overset{\leq k}{\sim} G_2$, /iii/ for every connected graph $H, |H| \leq k$, it is $\text{frq}(H, G_1) = \text{frq}(H, G_2)$ Now, for the case of equivalences we get:

Theorem 0.5. Let k be an integer, E_1, E_2 be equivalences. Following three properties are equivalent /i/ $E_1 \overset{k}{\sim} E_2$, /ii/ $E_1 \overset{\leq k}{\sim} E_2$, /iii/ for every $j \leq k$ $\text{frq}(K_j, E_1) = \text{frq}(K_j, E_2)$.

Proof. Use Remark 0.3., Theorem 0.4. and the fact, that only complete graphs are connected equivalences.

1. Frequencies in equivalences

Throughout this part of paper let A, B, C be equivalences, $A = s.K_u$, $B = K_v$ where $s > 0, u > 0, v > 0, Q = s.K_u + K_v$ and for every $i \leq u+v$ $Q_i = (s-1).K_u + K_i$.

Definition 1.1. We define two numbers for any equivalence E $\langle A, B \rangle \downarrow E = \text{card} \{Y, Z; E/Y \simeq A, E/Z \simeq B\} = \text{frq}(A, E) \cdot \text{frq}(B, E)$, $\langle A, B \rangle \downarrow \downarrow E = \text{card} \{Y, Z; E/Y \simeq A, E/Z \simeq B, Y \cup Z = X\}$, where $E = \langle X, R \rangle$.

Lemma 1.2. Let E be an equivalence. Then $\text{card} \{Y, Z; C/Y \simeq A, C/Z \simeq B, C/(Y \cup Z) \simeq E\} = [\langle A, B \rangle \downarrow \downarrow E] \cdot \text{frq}(E, C)$.

Proof. The number of the sets W such that $C/W \simeq E$ is $\text{frq}(E, C)$. For each such a set W we have $\langle A, B \rangle \downarrow \downarrow E$ ordered pairs $\langle Y, Z \rangle$ satisfying $C/Y \simeq A, C/Z \simeq B, W = Y \cup Z$.

Remark 1.3. Equivalences A, B in Lemma 1.2. can be arbitrary.

Lemma 1.4. The following equality is true

$$\langle A, B \rangle \downarrow \downarrow C = \sum_{i=1}^{u+v-1} [\langle A, B \rangle \downarrow \downarrow Q_i] \cdot \text{frq}(Q_i, C) + [\langle A, B \rangle \downarrow \downarrow Q_{u+v}] \cdot \text{frq}(Q_{u+v}, C) + [\langle A, B \rangle \downarrow \downarrow Q] \cdot \text{frq}(Q, C)$$

Proof. We denote $M_0 = \{ \langle Y, Z \rangle; C/Y \simeq A, C/Z \simeq B \}$ and further for every $i \leq u+v$ $M_i = \{ \langle Y, Z \rangle; C/Y \simeq A, C/Z \simeq B, C/(Y \cup Z) \simeq Q_i \}$. Finally, $M = \{ \langle Y, Z \rangle; C/Y \simeq A, C/Z \simeq B, C/(Y \cup Z) \simeq Q \}$. We have the disjoint decomposition $M_0 = M_1 \cup \dots \cup M_{u+v-1} \cup M_{u+v} \cup M$ and we can write $\text{card } M_0 = \sum_{i=1}^{u+v-1} \text{card } M_i + \text{card } M_{u+v} + \text{card } M$. Using Lemma 1.2. we obtain the

needed equality.

Lemma 1.5. Let $j+1 = s.u+v$, let E_1, E_2 be two equivalences such that $E_1 \overset{j}{\sim} E_2$ and $\text{frq}(Q, E_1) = \text{frq}(Q, E_2)$. Then $\text{frq}(Q_{u+v}, E_1) = \text{frq}(Q_{u+v}, E_2)$.

Proof. For $i \leq u+v-1$ it is $|Q_i| = (s-1).u+i \leq s.u+v-1 = j$ and by Theorem 0.5. $\text{frq}(Q_i, E_1) = \text{frq}(Q_i, E_2)$. Analogously, since $v \leq j$ and $s.u \leq j$ we have $\text{frq}(K_v, E_1) = \text{frq}(K_v, E_2)$ and $\text{frq}(s.K_u, E_1) = \text{frq}(s.K_u, E_2)$, i.e. $\langle s.K_u, K_v \rangle \downarrow E_1 = \langle s.K_u, K_v \rangle \downarrow E_2$. Now, we calculate using Lemma 1.4. $0 = \langle s.K_u, K_v \rangle \downarrow E_1 - \langle s.K_u, K_v \rangle \downarrow E_2 = \sum_{i=1}^{u+v-1} \langle s.K_u, K_v \rangle \downarrow \downarrow Q_i \cdot [\text{frq}(Q_i, E_1) - \text{frq}(Q_i, E_2)] + \langle s.K_u, K_v \rangle \downarrow \downarrow Q_{u+v} \cdot [\text{frq}(Q_{u+v}, E_1) - \text{frq}(Q_{u+v}, E_2)] + \langle s.K_u, K_v \rangle \downarrow \downarrow Q \cdot [\text{frq}(Q, E_1) - \text{frq}(Q, E_2)] = \langle s.K_u, K_v \rangle \downarrow \downarrow Q_{u+v} \cdot [\text{frq}(Q_{u+v}, E_1) - \text{frq}(Q_{u+v}, E_2)]$. Since $\langle s.K_u, K_v \rangle \downarrow \downarrow Q_{u+v} > 0$ we get finally $\text{frq}(K_{u+v}, E_1) = \text{frq}(K_{u+v}, E_2)$.

Definition 1.6. An equivalence E is called pseudoregular if there exist numbers $s \geq 0, u > 0, v > 0$ such that $E \simeq s.K_u + K_v$.

Now, we are able to prove the main theorem.

Theorem 1.7. Let k be an integer, E_1, E_2 be equivalences. Following four properties are equivalent /i/ $E_1 \overset{k}{\sim} E_2$, /ii/ $E_1 \overset{\leq k}{\sim} E_2$, /iii/ for every $j \leq k$ $\text{frq}(K_j, E_1) = \text{frq}(K_j, E_2)$, /iv/ for every $j \leq k$ there is a pseudoregular equivalence S_j such that $|S_j| = j$ and $\text{frq}(S_j, E_1) = \text{frq}(S_j, E_2)$.

Proof. To prove the theorem it suffices to show that the implication /iv/ \Rightarrow /iii/ is true. We use an indirect argument. If the implication is false there exist $i \leq k$ such that $\text{frq}(K_i, E_1) \neq \text{frq}(K_i, E_2)$. Let $i^\# = \min \{i; \text{frq}(K_i, E_1) \neq \text{frq}(K_i, E_2)\}$. Obviously $1 < i^\# \leq k$ and for $j = i^\# - 1$ we have by Theorem 0.5. $E_1 \overset{j}{\sim} E_2$. We know that $\text{frq}(S_{j+1}, E_1) = \text{frq}(S_{j+1}, E_2)$. Let $c = \min \{cp S; S \text{ is pseudoregular, } |S| = j+1, \text{frq}(S, E_1) = \text{frq}(S, E_2)\}$. Then $1 < c \leq cp S_{i^\#}$. Take $Q = s.K_u + K_v$ such that $s > 0, u > 0, v > 0, |Q| = j+1, cp Q = c = s+1$. By Lemma 1.5. $\text{frq}((s-1).K_u + K_{u+v}, E_1) = \text{frq}((s-1).K_u + K_{u+v}, E_2)$ contradicting the minimality of c because $op [(s-1).K_u + K_{u+v}] = s < c$.

Theorem 1.8. Let $k > 0, E_1, E_2$ be equivalences, $E_1 \overset{k}{\sim} E_2$. If there exists a pseudoregular equivalence S such that $|S| \leq k$ and $\text{frq}(S, E_1) = \text{frq}(S, E_2) = 0$ then $E_1 \overset{\sim}{\sim} E_2$.

Proof. Let $n = |E_1| = |E_2|$, let $S = s.K_u + K_v, |S| \leq k, \text{frq}(S, E_1) = \text{frq}(S, E_2)$. For every integer w define $S_w = s.K_u + K_w$. For $w \leq v$ we have $\text{frq}(S_w, E_1) = \text{frq}(S_w, E_2) = 0$ and by Theorem 1.7. /property /iv// $E_1 \overset{\sim}{\sim} E_2$. It is $1 = \text{frq}(E_1, E_1) = \text{frq}(E_1, E_2)$ and clearly $E_1 \overset{\sim}{\sim} E_2$.

2. Bounds of reconstructibility and nonreconstructibility

We are interested in the problem: for given k find n satisfying the implication $(|E_1| = |E_2| = n \text{ et } E_1 \stackrel{k}{\sim} E_2) \Rightarrow (E_1 \simeq E_2)$ where E_1, E_2 are arbitrary equivalences.

We denote $cp_i E$ the number of components of the equivalence E having at least i elements. Let us indicate two elementary facts: /fact 1/ $|E| = \sum_{i \geq 1} cp_i E$, /fact 2/ if $\text{frq}(s, K_1, E) \geq 1$ then $cp_i E \geq s$.

Theorem 2.1. Let $k > 2$, E_1, E_2 be equivalences, $|E_1| = |E_2| \leq k \cdot \ln(k/2)$ where \ln denotes the logarithmus naturalis. If $E_1 \stackrel{k}{\sim} E_2$ then $E_1 \simeq E_2$.

Proof. Suppose $E_1 \not\simeq E_2$ and define for every $i \leq k$ the integral part of k/i denoted $t_i = [k/i]$. Now, for every $i \leq k$ we have $\text{frq}(t_i, K_1, E_1) \geq 1$ by Theorem 1.8. and moreover $cp_i E_1 \geq t_i$ by /fact 2/.

We calculate $n = |E_1| = \sum_{i \geq 1} cp_i E_1 \geq \sum_{i=1}^k t_i \geq \sum_{i=1}^k (k/i - 1) = (k \cdot \sum_{i=1}^{k-1} 1/i) + 1 > k \cdot \ln(k/2) + 1$. We get a contradiction with the assumption that $n \leq k \cdot \ln(k/2)$.

Construction 2.2. For every $k \geq 1$ we construct two equivalences E_1, E_2 such that $E_1 \stackrel{k}{\sim} E_2$, $E_1 \not\simeq E_2$, $|E_1| = |E_2| = (k+1) \cdot 2^{k-1}$.

Proof. For $i=1, \dots, k+1$ we define the numbers a_i, b_i

$$a_i = \binom{n+1}{i} \text{ if } i \text{ is even} \quad b_i = 0 \text{ if } i \text{ is even}$$

$$0 \text{ if } i \text{ is odd} \quad \binom{n+1}{i} \text{ if } i \text{ is odd}.$$

The numbers a_i, b_i satisfy $a_i - b_i = (-1)^i \binom{n+1}{i}$, $a_i + b_i = \binom{n+1}{i}$.

We define $E_1 = \sum_{i=1}^{k+1} a_i \cdot K_i$, $E_2 = \sum_{i=1}^{k+1} b_i \cdot K_i$. It is obvious that $E_1 \not\simeq E_2$ because E_2 has 1-point components but E_1 has not. For every j ,

$1 \leq j \leq k$ we calculate $\text{frq}(K_j, E_1) - \text{frq}(K_j, E_2) = \sum_{i=j}^{k+1} a_i \cdot \binom{i}{j} - \sum_{i=j}^{k+1} b_i \cdot \binom{i}{j}$

$$= \sum_{i=j}^{k+1} (a_i - b_i) \cdot \binom{i}{j} = \sum_{i=j}^{k+1} (-1)^i \cdot \binom{k+1}{i} \cdot \binom{i}{j} = 0$$

and we get $\text{frq}(K_j, E_1) = \text{frq}(K_j, E_2)$. It is $E_1 \stackrel{k}{\sim} E_2$ by Theorem 1.7.

Finally, we calculate $|E_1| + |E_2| = \sum_{i=1}^{k+1} a_i \cdot i + \sum_{i=1}^{k+1} b_i \cdot i =$

$$= \sum_{i=1}^{k+1} (a_i + b_i) \cdot i = \sum_{i=1}^{k+1} \binom{k+1}{i} \cdot i = (k+1) \cdot 2^k$$

which yields $|E_1| = |E_2| = (k+1) \cdot 2^{k-1}$.

Remark 2.3. In [6] we have defined reconstructibility indicating function $u_{\mathcal{E}}$ of the class of graphs \mathcal{E} . If we denote \mathcal{E} the class of all equivalences we can write the result of this paper in the form: for every $k > 2$ $k \cdot \ln(k/2) \leq u_{\mathcal{E}}(k) < (k+1) \cdot 2^{k-1}$.

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