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PARTIALLY DIRECTED MOORE GRAPHS

JURAJ BOSÁK

1. Introduction

Undirected T-graphs (sometimes called strongly geodetic graphs) have been studied in several papers [3, 7, 13]. Plesník and Znáám [14] have dealt with analogous problems for directed graphs. In the present paper we generalize this concept for partially directed graphs. We show that any T-graph is either an undirected tree or a homogeneous graph with a finite diameter, i.e., a tied graph. The most interesting T-graphs — called Moore graphs — are finite tied graphs. We get a class much richer than the class of undirected Moore graphs (cf. [1, 2, 8, 10, 12, 13]). For example, there exist infinitely many (partially directed) Moore graphs of diameter two. However, given the number of vertices of a Moore graph G of diameter two the following invariants are uniquely determined: the valency of G , the directed and the undirected valencies of G , the degree of G and the spectrum of G (that is, the eigenvalues of the adjacency matrix of G and their multiplicities). Therefore our main attention is devoted to finding conditions for the existence of a Moore graph of diameter two with a given number of vertices. Some results are concerned with the Moore graphs of other diameters, with the T-graphs or a generalization of them, namely the P-graphs. Several open problems are stated.

Most results of the present paper were reported at international conferences held in Kalamazoo [4] and Keszthely [5] in 1976.

2. Notation and terminology

We shall give here only definitions of the notions that are not clear from the context.

Graphs considered in this paper are *partially directed*, i.e., they may contain directed edges as well as undirected ones. A graph is called *directed* [*undirected*] if all its edges are directed [undirected, respectively]. A graph is called *mixed* if it

contains at least one directed edge and at least one undirected edge. An edge is said to be a *link* [a *loop*] if its end vertices are different [equal] (cf. [15]). Two edges with the same end vertices are called *parallel*. If they are both undirected or both directed with the same orientation, they are called *multiple*. Two directed parallel links that are not multiple form a *pair of oppositely directed links*. A graph is said to be *complete* if it has no loops and if any two different vertices are joined either by exactly one undirected link or by just a pair of oppositely directed links. The undirected complete graph with n vertices will be denoted by K_n . A graph is said to be *finite* [*infinite*] if it has a finite [an infinite] number of elements (vertices and edges). By a *factor* of a graph G we mean a subgraph of G containing all the vertices of G . Graphs G and H are said to be *isomorphic* if there exists a bijection from the vertex set of G onto the vertex set of H preserving the number (generally, the cardinal number) of both directed (in each orientation) and undirected edges joining any two vertices.

Let v be a vertex of a graph G . Denote by E_1 the (cardinal) number of undirected links incident with v , by E_2 [E_3] the number of links directed from [to] v and by E_4 [E_5] the number of directed [undirected] loops incident with v . Then the following cardinal numbers can be defined:

$$\begin{aligned} \deg_G v &= E_1 + E_2 + E_3 + E_4 + 2E_5 && \text{(the degree of } v \text{ in } G), \\ \text{dd}_G v &= E_2 + E_3 + E_4 && \text{(the directed degree of } v \text{ in } G), \\ \text{ud}_G v &= E_1 + 2E_5 && \text{(the undirected degree of } v \text{ in } G), \\ \text{od}_G v &= E_1 + E_2 + E_4 + E_5 && \text{(the outdegree of } v \text{ in } G), \\ \text{dod}_G v &= E_2 + E_4 && \text{(the directed outdegree of } v \text{ in } G), \\ \text{id}_G v &= E_1 + E_3 + E_4 + E_5 && \text{(the indegree of } v \text{ in } G), \\ \text{did}_G v &= E_3 + E_4 && \text{(the directed indegree of } v \text{ in } G), \\ \text{val}_G v &= E_1 + \max \{E_2, E_3\} + E_4 + E_5 && \text{(the valency of } v \text{ in } G), \\ \text{dv}_G v &= \max \{E_2, E_3\} + E_4 && \text{(the directed valency of } v \text{ in } G), \\ \text{uv}_G v &= E_1 + E_5 && \text{(the undirected valency of } v \text{ in } G). \end{aligned}$$

The *degree* [*valency*, *directed valency*, *undirected valency*] of a graph G is defined as the supremum of the set of the degrees [valencies, directed valencies, undirected valencies, respectively] of all the vertices of G .

Let d be a cardinal number. A graph G is said to be *regular of degree* d if $\deg_G v = d$ for every vertex v of G . A graph G is said to be *homogeneous of valency* d if $\text{od}_G v = \text{id}_G v = d$ for every vertex v of G . In the case of G being an undirected loopless graph these two concepts coincide.

By a *semitrail* from u to v in a graph G we mean a finite sequence

$$S = [v_0, e_1, v_1, e_2, v_2, \dots, v_{n-1}, e_n, v_n],$$

where n is a non-negative integer (called the length of S); $v_0 = u, v_1, v_2, \dots, v_{n-1}, v_n = v$ are vertices of G ; e_1, e_2, \dots, e_n are mutually different edges of G and

(provided that $n \geq 1$) for $i \in \{1, 2, \dots, n\}$ v_{i-1} and v_i are the end vertices of e_i . If, moreover, every e_i is either undirected, or directed from v_{i-1} to v_i , then S is said to be a *trail*. A semitrail [trail] whose vertices are mutually different, is called a *semipath* [path]. A semitrail [trail] from u to v is said to be a *semicircuit* [circuit] if it has a positive length and if its vertices are mutually different with the exception of $u = v$. The undirected graph consisting of the vertices and the edges of a path with n vertices will be denoted by P_n . The graph consisting of the vertices and the edges of a circuit will be called a *cycle*. A directed [undirected] cycle with n edges will be denoted by $Z_n[C_n]$.

A graph G is said to be a *tree* if for each ordered pair $[u, v]$ of vertices of G there exists in G exactly one semitrail from u to v .

Let u and v be vertices of a graph G . Denote by $\rho_G(u, v)$ the smallest length of a path from u to v . If there is no such path, we put $\rho_G(u, v) = \infty$. A graph G is said to be *connected* [strongly connected] if for every ordered pair $[u, v]$ of vertices of G there exists a semipath [path] from u to v . The *diameter* of G is defined as $k_G = \sup \rho_G(u, v)$, where the supremum is taken through all the ordered pairs $[u, v]$ of vertices of G . If G is not strongly connected, then $k_G = \infty$.

A vertex v of a graph G is said to be a *cutpoint* of G if there exist two different edges e and f of G such that every semitrail containing e and f contains v between e and f . A maximal connected subgraph H of G containing no cutpoint of H is called a *block* of G .

Let p be a positive integer and G be a finite graph with p vertices v_1, v_2, \dots, v_p . By the *adjacency matrix* of G we mean the $p \times p$ matrix $A = [a_{i,j}]$, where $a_{i,j}$ is the number of edges of G joining v_i and v_j that are either undirected or directed from v_i to v_j ($i, j \in \{1, 2, \dots, p\}$).

3. T-graphs

A (partially directed) graph G is said to be a *T-graph* if for each ordered pair $[u, v]$ of vertices of G there exists in G exactly one trail from u to v of a length not greater than the diameter of G . (In [3, 7, 13, 14] T-graphs were called strongly geodetic graphs.) In [7] the following result has been proved:

Lemma 1. *Every undirected T-graph is either a tree or a regular graph with a finite diameter.*

A directed version of this result has been proved in [14]. We shall generalize it for partially directed graphs.

Theorem 1. *Every T-graph is either an undirected tree or a homogeneous graph with a finite diameter.*

Proof. Let G be a T-graph of diameter k . Distinguish three cases:

I. The diameter k is infinite. Let there exist in G an edge e directed from a vertex u to a vertex v . If $v = u$, then there are in G two trails from u to u , namely $[u, e, u]$ and $[u]$. If $v \neq u$, then there is in G a trail from v to u ,

$$[v, e_1, v_1, e_2, v_2, \dots, u]$$

so that there are again two trails from u to u , namely

$$[u, e, v, e_1, v_1, e_2, v_2, \dots, u]$$

and $[u]$. As these possibilities lead to a contradiction, we may suppose that all edges of G are undirected. But then the assertion follows from Lemma 1.

II. The diameter $k = 0$ or $k = 1$. Then every T-graph is homogeneous and the assertion holds.

III. The diameter k of G is finite, $k \geq 2$. According to Lemma 1 we may suppose that in G there exists a directed edge. We shall prove that every directed edge of G belongs to a directed cycle with $k + 1$ edges. Let e be an edge directed from u to v (evidently $v \neq u$). Then there exists a trail from v to u

$$[v, e_1, v_1, \dots, u]$$

of a length $\leq k$ so that

$$C = [u, e, v, e_1, v_1, \dots, u]$$

is a circuit of a length $\leq k + 1$. Obviously C cannot contain an undirected edge as then its end vertices would be connected by two trails of lengths $\leq k$. As G is a T-graph, G cannot have a directed circuit of a length $\leq k$. Hence C has the length $k + 1$. Thus e lies in a directed cycle consisting of elements of C and this cycle has $k + 1$ edges.

Now we shall prove that if x is any vertex of G , then $\text{id}_G x = \text{od}_G x$. Since undirected edges have no influence upon the validity of this equality, we may restrict ourselves to directed edges. (In fact, we shall prove that $\text{did}_G x = \text{dod}_G x$.) Every edge directed from x , as we have proved, lies in a directed cycle with $k + 1$ edges. As G is a T-graph, no edge can be contained in two such cycles so that they are edge-disjoint. In each of the cycles there is an edge directed to x so that $\text{od}_G x \leq \text{id}_G x$. Analogously the inequality $\text{id}_G x \leq \text{od}_G x$ can be proved. Hence $\text{id}_G x = \text{od}_G x$.

Let y and z be vertices of G such that $\rho_G(y, z) = k$. We shall prove that $\text{od}_G y = \text{id}_G z$. Let

$$[y, e_1, v_1, e_2, v_2, \dots, v_{k-1}, e_k, z]$$

be a path of the length k from y to z . Let y_1, y_2, \dots, y_s be all the vertices of G different from v_1 such that there exists an edge directed from y to one of them or an undirected edge joining y to one of them. Evidently for $i \in \{1, 2, \dots, s\}$ we have

$\rho_G(y_i, z) = k$ and no two of the corresponding paths of length k have a vertex different from z in common. Therefore $\text{od}_G y \leq \text{id}_G z$. Analogously, considering the edges directed to z , the inequality $\text{od}_G y \geq \text{id}_G z$ can be obtained. Consequently, $\text{od}_G y = \text{id}_G z$.

Let e be an edge of G directed from a vertex U to a vertex V . Then evidently $\rho_G(V, U) = k$ so that according to what has been just proved $\text{id}_G V = \text{od}_G V = \text{id}_G U = \text{od}_G U$. The edge e lies in a circuit $C(e)$ of the length $k + 1$ with all edges directed. Let the vertices of $C(e)$ be $v_0 = U, v_1 = V, v_2, \dots, v_k$ such that there exists an edge directed from v_{i-1} to v_i ($i = 1, 2, \dots, k$) and an edge directed from v_k to v_0 . Using the preceding results we get

$$\text{id}_G v_i = \text{od}_G v_i = \text{id}_G V$$

for each $i \in \{0, 1, \dots, k\}$.

Let w be a vertex of G not contained in $C(e)$. Let v_j be the vertex of $C(e)$ such that

$$\rho_G(w, v_j) \leq \rho_G(w, v_i)$$

for all $i = 0, 1, 2, \dots, k$. Put $\rho_G(w, v_j) = J, 1 \leq J \leq k$. Then the vertex v_{j+k-J} has the property that

$$\rho_G(w, v_{j+k-J}) = k,$$

where the index $j + k - J$ is to be taken modulo $k + 1$. From the preceding parts of the proof it follows that

$$\text{id}_G w = \text{od}_G w = \text{id}_G V.$$

Thus G is a homogeneous graph. Q.E.D.

A homogeneous T -graph of valency d and with a finite diameter k will be called a tied graph of type (d, k) . Thus from Theorem 1 it follows:

Corollary 1. *A graph G is a T -graph if and only if G is either an undirected tree or a tied graph.*

Evidently, we have:

Lemma 2. *A tied graph of type (d, k) is finite if and only if its valency d is finite.*

Concerning infinite tied graphs, the following result is known:

Proposition 1. *For every infinite cardinal number d and finite cardinal number k there exists an undirected [directed] tied graph of type (d, k) .*

The undirected version of this result has been proved in [3, 7, 11], the directed one in [14]. Using the method of these papers it is easy to prove that the following version of Proposition 1 is true:

Proposition 2. *For every infinite cardinal number d and a finite cardinal number k there exists a mixed tied graph of type (d, k) .*

In view of these results we shall be concerned now only with finite tied graphs. A finite tied graph of type (d, k) will be called also a Moore graph of type (d, k) . (For undirected Moore graphs see [1, 2, 8, 10, 12, 13].)

4. Moore graphs of diameter two

Theorem 2. *Let G be a Moore graph of diameter two. Then there exist non-negative integers z and r satisfying the following conditions:*

- A. $\text{dod}_G v = \text{did}_G v = z$ for every vertex v of G .
- B. $\text{uv}_G v = r$ for every vertex v of G .
- C. The number of vertices of G is

$$(1) \quad p = (z + r)^2 + z + 1 \geq 3$$

and exactly one of the following cases occurs:

- I. $z = 1, r = 0, p = 3$.
- II. $z = 0, r = 2, p = 5$.
- III. There is an odd positive integer c such that

$$(2) \quad c \mid (4z - 3)(4z + 5)$$

and

$$(3) \quad r = \frac{1}{4}(c^2 + 3).$$

Proof. G is a homogeneous graph. Therefore there exists a non-negative integer d such that

$$\text{val}_G v = \text{od}_G v = \text{id}_G v = d$$

for every vertex v of G . Thus

$$\text{dv}_G v = \text{dod}_G v = \text{did}_G v$$

for every vertex v of G . Put

$$\begin{aligned} z(v) &= \text{dv}_G v, \\ r(v) &= d - z(v). \end{aligned}$$

Using trails of lengths 0, or 1, or 2 beginning with v we can arrive at 1, or d , or $r(v)(d - 1) + z(v)d$ vertices of G , respectively. Since G is a T -graph of diameter two, we have

$$p = 1 + d + r(v)(d - 1) + z(v)d$$

so that

$$(4) \quad p = 1 + d + d^2 - r(v).$$

As p and d do not depend on v , the same is valid for

$$r(v) = d^2 + d + 1 - p$$

and

$$z(v) = d - r(v).$$

Put $r(v) = r$, $z(v) = z$. Thus A and B hold. Moreover, substituting $r(v) = r$ and

$$(5) \quad d = z + r$$

in (4) we obtain (1) (evidently $p \geq 3$ as G has diameter 2).

Let A be the adjacency matrix of G . Then

$$A^2 + A - (r - 1)I = J,$$

where I is the identity matrix of order p and J is the matrix of order p each entry of which is 1. The eigenvalues of J are

$$\lambda_1 = \lambda_2 = \dots = \lambda_{p-1} = 0, \quad \lambda_p = p.$$

Thus A has the eigenvalues $\mu_1, \mu_2, \dots, \mu_p$ such that $\mu_1, \mu_2, \dots, \mu_{p-1}$ satisfy the equation

$$(6) \quad x^2 + x - (r - 1) = 0$$

and μ_p satisfies the equation

$$x^2 + x - (r - 1) = p.$$

As G is a homogeneous graph of valency $d = z + r$, A has an eigenvalue d . From (1) it follows that $d = \mu_p$. The roots of (6) are

$$(7) \quad s = \frac{-1 + \sqrt{4r - 3}}{2}, \quad t = \frac{-1 - \sqrt{4r - 3}}{2}.$$

(For $r = 0$ we put $\sqrt{4r - 3} = i\sqrt{3}$.) Thus A has at most 3 different eigenvalues, namely d , s and t . The eigenvalue d is, according to the preceding considerations, simple (cf. [16]). Denote the multiplicity of s or t in A by a or b , respectively. Thus

$$(8) \quad a + b + 1 = p.$$

As the trace of A is zero, the sum of the eigenvalues of A is

$$(9) \quad as + bt + d = 0.$$

Substituting (7) into (9) gives

$$(10) \quad a + b + \sqrt{4r - 3}(b - a) = 2d.$$

Distinguish two cases:

Case 1. $b = a$. Then from (10) and (8) it follows that $2d = a + b = p - 1$ so that

$$(11) \quad p = 2d + 1.$$

Since $p \geq 3$, we have

$$(12) \quad d \geq 1.$$

Using (11), (1) and (5) we get

$$2d + 1 = p = (z + r)^2 + z + 1 = d^2 + d - r + 1$$

so that

$$(13) \quad r = d^2 - d.$$

However, $r \leq d$. Hence $d^2 - d \leq d$ and (12) implies

$$(14) \quad 1 \leq d \leq 2.$$

According to (14), (13), (5) and (1) we have only two possibilities:

- I. $d = 1, r = 0, z = 1, p = 3$.
- II. $d = 2, r = 2, z = 0, p = 5$.

Case 2. $b \neq a$. Put

$$(15) \quad c = \sqrt{4r - 3}.$$

Then (3) holds. From (10) it follows that $r \neq 0$, hence $r \geq 1$ and $c \geq 0$. By (10) c is rational. As $4r - 3$ is an integer, c is an integer, too. From (10), (8) and (15) it follows that

$$(16) \quad p - 1 - 2d + c(p - 2a - 1) = 0.$$

If in (16) we successively substitute

$$\begin{aligned} p &= d^2 + d - (r - 1), \\ d &= r + z, \\ r &= \frac{1}{4}(c^2 + 3), \end{aligned}$$

we get after simple calculations

$$\begin{aligned} c^5 + c^4 + c^3(8z + 6) + c^2(8z - 2) + c(16z^2 + 40z - 32a + 9) + \\ + (16z^2 + 8z - 15) = 0. \end{aligned}$$

Since

$$16z^2 + 8z - 15 = (4z - 3)(4z + 5)$$

and c is an integer, we get (2). Consequently, c is an odd integer. Q.E.D.

Remark. If we put in Theorem 2 $z = 0$ or $r = 0$, respectively, we easily obtain the following two well-known results.

Corollary 2 [10]. Any undirected Moore graph of diameter two has 5, 10, 50 or 3250 vertices.

Corollary 3 [14]. A directed graph is a Moore graph of diameter two if and only if it is isomorphic to Z_3 .

Theorem 3. Let there exist a Moore graph G of diameter two with p vertices. Then the following invariants d, z, r, D, s, t, a, b of G are uniquely determined by p :

$$(17) \quad d = [\sqrt{p}] \quad (\text{the valency of } G),$$

$$(18) \quad z = p - d^2 - 1 \quad (\text{the directed valency of } G),$$

$$(19) \quad r = d - z \quad (\text{the undirected valency of } G),$$

$$(20) \quad D = 2z + r \quad (\text{the degree of } G),$$

$$(21) \quad s = \frac{-1+c}{2},$$

$$(22) \quad t = \frac{-1-c}{2} \quad (\text{the eigenvalues of the adjacency matrix of } G \text{ different from } d),$$

$$(23) \quad a = \frac{\frac{1}{2}(c+1)(p-1) - d}{c} \quad (\text{the multiplicity of } s),$$

$$(24) \quad b = p - 1 - a \quad (\text{the multiplicity of } t),$$

where

$$(25) \quad c = \sqrt{4r-3}.$$

(We put $\sqrt{-3} = i\sqrt{3}$.)

Proof. Denote the valency, directed valency, undirected valency and the degree of G by d, z, r and D , respectively. Evidently, (19) and (20) hold. According to (1) we have

$$p = d^2 + z + 1.$$

Therefore we get (18) and the inequalities

$$d^2 < d^2 + z + 1 \leq d^2 + d + 1 < d^2 + 2d + 1 = (d+1)^2$$

imply

$$d < \sqrt{p} < d+1.$$

Hence we obtain (17). Define c by (25). Then (21) and (22) follow from (7). Solving the system of the two equations (8) and (10) we get (23) and (24). Q.E.D.

Theorem 4. *Let there exist a Moore graph of diameter two with p vertices. Then either $p \in \{3, 5\}$ or p is an even integer of the form*

$$(26) \quad p = d(d+1) - s(s+1) \geq 6,$$

where d and s are integers, $d \geq 2$ (d is the valency of the graph and the greatest eigenvalue of its adjacency matrix), $s \geq 0$ (s is the second greatest eigenvalue),

$$(27) \quad s(s+1) < d$$

and

$$(28) \quad 2s+1 \mid (2d+1)(2d-3).$$

Proof. We use the same notation as in the preceding theorems. Thus p, d, z and r are integers, $p \geq 3$, $0 \leq r \leq d$, $0 \leq z \leq d$, $d = z + r$. According to (1) we have

$$p = d^2 + d - (r-1).$$

From (3) and (21) it follows that

$$(29) \quad r-1 = \frac{1}{4}(c^2+3) - 1 = s(s+1)$$

and we get (26).

According to Theorem 2 either $p \in \{3, 5\}$ or c is an odd positive integer and by (21) s is a non-negative integer. But then from (26) we get that p is an even integer. Evidently, $p \neq 2$. If $p = 4$, then (17) and (18) give $d = 2$, $z = -1$, a contradiction. Therefore $p \geq 6$. Now (17) implies $d \geq 2$. From (29) we get

$$s(s+1) = r-1 < r \leq d$$

and (27) holds. It remains to prove (28). We shall use (2). From (21) we get

$$c = 2s + 1.$$

As $z = d - r = d - 1 - s(s+1)$, we have

$$(4z-3)(4z+5) = (4z+1)^2 - 16 = (4d-3-4s(s+1))^2 - 16.$$

However,

$$4s(s+1) = (2s+1)^2 - 1 \equiv -1 \pmod{2s+1}.$$

Thus

$$\begin{aligned} (4z-3)(4z+5) &\equiv (4d-3+1)^2 - 16 = \\ &= 4(2d+1)(2d-3) \pmod{2s+1}. \end{aligned}$$

Therefore (2) is equivalent to

$$2s+1 \mid 4(2d+1)(2d-3).$$

Since $2s+1$ is relatively prime to 4, we get (28). Q.E.D.

Remark. If $s = 0$, the conditions (27) and (28) always hold. Checking (28) can be simplified in the case when $2s + 1$ is a prime power (i.e., if $s \in \{1, 2, 3, 5, 6, 8, 9, 11, 14, \dots\}$):

Proposition 3. *Let d and s be integers and let $2s + 1$ be a power of a prime. Then (28) holds if and only if*

$$d \equiv s \quad \text{or} \quad s + 2 \pmod{2s + 1}.$$

Proof. (28) can be written in the form

$$(2d + 1)(2d - 3) \equiv 0 \pmod{2s + 1}.$$

Let $2s + 1$ be a power of a prime p . Evidently, $p \neq 2$. As

$$(2d + 1) - (2d - 3) = 4$$

is not divisible by p , we have either

$$2d + 1 \equiv 0 \pmod{2s + 1},$$

or

$$2d - 3 \equiv 0 \pmod{2s + 1}.$$

The former congruence has the unique solution $d \equiv s$, the latter $d \equiv s + 2 \pmod{2s + 1}$. Q.E.D.

We need the following result proved in [10] and closely related to Corollary 2:

Lemma 3. *Undirected Moore graphs of diameter two with p vertices for $p \in \{5, 10, 50\}$ exist and are unique up to isomorphism.*

The Moore graph of diameter two with 5 vertices is C_5 (pentagon), with 10 vertices is the Petersen graph, with 50 is the Hoffman-Singleton graph (see [10, 2]).

Theorem 5. *A Moore graph of diameter two exists for every $p \in \{3, 5, 10, 18, 50\}$ and every $p = d(d + 1)$, where d is an integer, $d \geq 2$.*

Proof. For $p = 3$ it is sufficient to take the directed cycle Z_3 (cf. Corollary 3). The cases $p \in \{5, 10, 50\}$ are included in Lemma 3.

Let $p = 18$. We construct a Moore graph M of diameter two with 18 vertices as follows: The vertices of M are the residue classes $1, 2, 3, \dots, 18$ modulo 18. An arbitrary vertex x is incident with one edge directed from x to a vertex

$$\text{dir } x = \begin{cases} x + 1, & \text{if } x \equiv 1 \text{ or } 2 \pmod{6}; \\ x - 2, & \text{if } x \equiv 3 \pmod{6}; \\ x + 2, & \text{if } x \equiv 4 \pmod{6}; \\ x - 1, & \text{if } x \equiv 5 \text{ or } 6 \pmod{6} \end{cases}$$

and with three undirected edges joining x with the vertices

$$x + 3, \quad x - 3, \quad \text{dir } x + 9.$$

The graph M (decomposed into two factors) is given in Fig. 1. (Another representation of M is given in [4].)

Finally, let $p = d(d + 1)$, $d \geq 2$. We shall construct a Moore graph $B(d)$ of type $(d, 2)$. (For $d = 3$ see Fig. 2.) The vertices of $B(d)$ are ordered pairs $[x, y]$ of

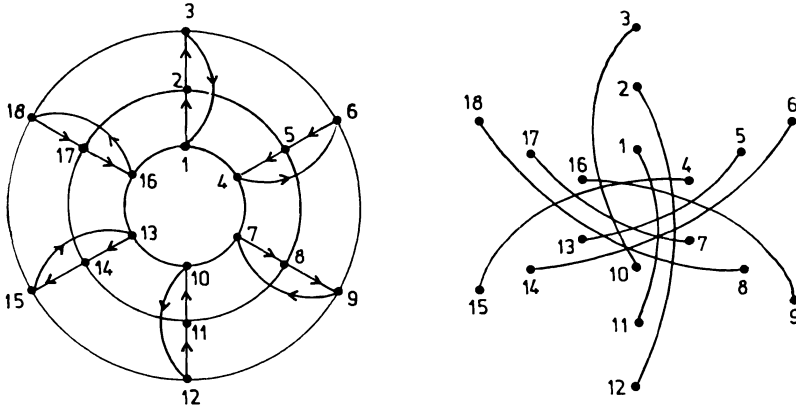


Fig. 1. A Moore graph M with 18 vertices

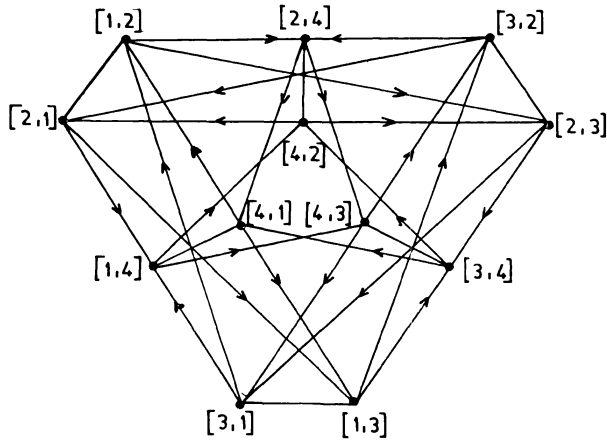


Fig. 2. A Moore graph $B(3)$ with 12 vertices

integers such that $1 \leq x \leq d + 1$, $1 \leq y \leq d + 1$, $x \neq y$. An arbitrary vertex $[x, y]$ is joined with a vertex $[X, Y]$ by an undirected edge if and only if $X = y$, $Y = x$ and by an edge directed from $[x, y]$ to $[X, Y]$ if and only if $X = y$, $Y \neq x$. (The graph $B(d)$ is isomorphic to the graph $B^+(d, 1)$ from [4].)

It is easy to check that all these graphs are Moore graphs of diameter two. Q.E.D.

Table 1

p	d	z	r	D	c	s	t	a	b	example
3	1	1	0	2	$i\sqrt{3}$	ϵ	$\bar{\epsilon}$	1	1	Z_3
5	2	0	2	2	$\sqrt{5}$	α	β	2	2	C_5
6	2	1	1	3	1	0	-1	3	2	$B(2)$
10	3	0	3	3	3	1	-2	5	4	Petersen graph
12	3	2	1	5	1	0	-1	8	3	$B(3)$
18	4	1	3	5	3	1	-2	10	7	M
20	4	3	1	7	1	0	-1	15	4	$B(4)$
30	5	4	1	9	1	0	-1	24	5	$B(5)$
40	6	3	3	9	3	1	-2	24	15	unknown
42	6	5	1	11	1	0	-1	35	6	$B(6)$
50	7	0	7	7	5	2	-3	28	21	Hoffman-Singleton graph
54	7	4	3	11	3	1	-2	33	20	unknown
56	7	6	1	13	1	0	-1	48	7	$B(7)$
72	8	7	1	15	1	0	-1	63	8	$B(8)$
84	9	2	7	11	5	2	-3	48	35	unknown
88	9	6	3	15	3	1	-2	55	32	unknown
90	9	8	1	17	1	0	-1	80	9	$B(9)$

A survey of “small” Moore graphs ($p \leq 100$) is given in Table 1, where

$$\epsilon = \frac{-1 + i\sqrt{3}}{2}, \quad \bar{\epsilon} = \frac{-1 - i\sqrt{3}}{2},$$

$$\alpha = \frac{-1 + \sqrt{5}}{2}, \quad \beta = \frac{-1 - \sqrt{5}}{2}.$$

Problem 1. Are there two non-isomorphic Moore graphs of diameter two with the same number of vertices?

Problem 2. Are there Moore graphs of diameter two with 40, 54, 84, or 88 vertices?

5. Moore graphs of diameter $k \neq 2$

The cases $k = 0$ or 1 are not interesting as we obviously have:

Lemma 4. Every finite graph of diameter 0 is a Moore graph.

Lemma 5. A finite graph G of diameter 1 is a Moore graph if and only if G is loopless and simple (so that G is a complete graph).

For diameters $k \geq 3$ the following results are known:

Lemma 6. [1, 8]. *An undirected graph of a finite diameter $k \geq 3$ is a Moore graph if and only if it is isomorphic to C_{2k+1} .*

Lemma 7 [14]. *A directed graph of a finite diameter $k \geq 3$ is a Moore graph if and only if it is isomorphic to Z_{k+1} .*

From Theorem 1, Corollary 3, Proposition 1, and Lemmas 2, 4, 5 and 7 it immediately follows:

Corollary 4 [14]. *A directed graph G is a T -graph if and only if one of the following cases occurs:*

- (i) G consists of a single vertex and an arbitrary number of loops.
- (ii) G is a complete directed graph.
- (iii) G is a directed cycle.
- (iv) G is an infinite directed tied graph.

Problem 3. *Is there a mixed Moore graph of diameter $k \geq 3$?*

6. P-graphs

A (partially directed) graph G is said to be a P -graph if for each ordered pair $[u, v]$ of vertices of G there exists in G exactly one path from u to v of a length not greater than the diameter of G .

A graph G is said to be a quasitree if for each ordered pair $[u, v]$ of vertices of G there exists exactly one path from u to v .

Theorem 6. *A graph G is a quasitree if and only if G is connected and every block of G is isomorphic to K_2 , C_1 , or a directed cycle.*

Proof. Let G be a connected graph whose every block is K_2 , C_1 , or a directed cycle. Then evidently G is a quasitree.

Conversely, let G be a quasitree. If we replace every undirected loop of G by a directed loop and every undirected link of G by a pair of oppositely directed links, we get a directed graph G^* that is again a quasitree. Therefore every edge of G^* is contained in exactly one (directed) circuit.

Evidently, it now suffices to prove that every block of G^* is a (directed) cycle. Let B be a block of G^* with n vertices. If $n = 1$, then B is isomorphic to Z_1 . Therefore, let $n \geq 2$. Then B has no loops. Let C be a cycle contained in B . Evidently, every edge joining two vertices of C lies in C . Suppose that there exists a vertex v of B that does not lie in C . As B is a block, there exists a semipath $[u_0, e_1, u_1, e_2, \dots, u_r = v, \dots, u_{s-1}, e_s, u_s]$, where $s \geq 2$, $1 \leq r \leq s - 1$, u_0 and u_s are in C , but u_1, u_2, \dots, u_{s-1} are not in C . Without loss of generality we may suppose that e_1 is directed from u_0 to u_1 . Each e_i , $i \in \{2, 3, \dots, t\}$ lies in a cycle of B . Therefore in B

there exists a path $[v_1 = u_1, f_1, v_2, \dots, v_{t-1}, f_t, v_t = u_s]$. But then in B there are two paths from u_0 to u_s , the first is inside C and the second one is

$$[u_0, e_1, u_1 = v_1, f_1, v_2, \dots, v_{t-1}, f_t, v_t = u_s].$$

However, this is a contradiction to the definition of a quasitree. Q.E.D.

Graphs G and H are said to be similar if deleting all the loops and replacing every undirected link by a pair of oppositely directed links in both G and H yields two isomorphic (directed) graphs. We obviously have:

- Proposition 4.** 1. Every graph similar to a T -graph is a P -graph.
2. Every quasitree is a P -graph.

Problem 4. Is there a P -graph that is neither a quasitree nor a graph similar to a T -graph?

Evidently, we have:

Lemma 8. Let G be a loopless undirected graph. Then G is a P -graph if and only if G is a T -graph.

From Lemmas 2—6, 8 and Corollaries 1—2 we immediately have:

Corollary 5 (cf. [7], Theorem 3). A loopless undirected graph G is a P -graph (or, equivalently, a T -graph) if and only if one of the following cases occurs:

- (i) G is a tree.
- (ii) G is a complete graph.
- (iii) G is a cycle with an odd number of edges.
- (iv) G is the Petersen graph.
- (v) G is the Hoffman-Singleton graph.
- (vi) G is a Moore graph of type $(57, 2)$.
- (vii) G is an infinite tied graph.

Remark 1. The existence and uniqueness of graphs (vi) from Corollary 5 is an open problem (see [10, 13]).

Remark 2. The condition "loopless" in Corollary 5 is not an essential restriction as a graph G is a P -graph if and only if deleting all loops from G yields a P -graph. On the other hand, a T -graph with a non-zero diameter cannot have loops.

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ЧАСТИЧНО ОРИЕНТИРОВАННЫЕ ГРАФЫ МУРА

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Резюме

(Частично ориентированный) граф G называется Т-графом [Р-графом] если для всякой упорядоченной пары $[u, v]$ его вершин существует в G точно одна цепь [точно один путь, соответственно] длины непревышающей диаметр графа G . Граф G называется однородным валентности d , если внешняя и внутренняя степень всякой вершины равны d . Показано, что всякий Т-граф является или неориентированным деревом, или однородным графом конечного диаметра — завязанным графом. Конечные завязанные графы называются графами Мура. В работе изучаются свойства Р-графов, Т-графов и графов Мура. Главные результаты касаются вопросов существования графов Мура диаметра два — кроме других результатов — построен бесконечный класс таких графов.