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THE REFINEMENT OF TWO ISOMORPHIC
GENERALIZED LEXICOGRAPHIC PRODUCTS

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INTRODUCTION

In a recent paper [3] M. Novotný has described a decomposition induced by an isomorphism of two Cartesian products. The purpose of this note is to introduce a generalization of the lexicographic product and to extend some of the results of [3] and [4] to this product.

1.

Without reference we shall use the terminology and notation of [3]. We start with a generalization of the lexicographic product. Let A , M be nonempty sets and let M^A denote the set of all mappings from A to M . The elements of M^A will be denoted also by $(m_a)_{a \in A}$. A subset ρ of M^A will be called an A -relation on M . If card $A = h < \aleph_0$, we agree to identify the A -relations on M with the h -ary relations on M , i.e. with the subsets of the Cartesian power $M^h = M \times \dots \times M$.

Let A , C , Q and $B_q (q \in Q)$ be nonempty sets. Let α_q be A -relations on B_q and let γ be a C -relation on Q . The Cartesian product of the sets B_q will be denoted by $B = \prod_{q \in Q} B_q$. We shall consider the set Φ of all logical formulae which:

1° have exactly the free variables $(f_q^a)_{q \in Q} \in B$ and

2° are built up from the following atomic predicates:

- (i) the equalities: $E_{a'a''} \equiv_{df} f_q^{a'} = f_q^{a''}$ ($q \in Q$, a' , $a'' \in A$),
- (ii) the predicates defined by the relations α_q : $P_{\nu q} \equiv_{df} (f_q^a)_{a \in A} \in \alpha_q$ ($q \in Q$, $\nu : A \rightarrow A$),
- (iii) the predicates defined by the relation γ : $W_\mu \equiv_{df} (\mu c)_{c \in C} \in \gamma$ ($\mu : C \rightarrow Q$). Hence the formulae from Φ are formed from the atomic predicates (i)—(iii) by means of the disjunction \vee , the conjunction $\&$, the negation \sim , and the quantifiers \exists and \forall (whose bound variables range over Q) according to the laws of the predicate calculus.

Example. Let card $A = \text{card } C = 2$ and let \leq_q and \leq be binary reflexive relations on $B_q (q \in Q)$ and Q respectively. We have the following example of a formula from Φ :

$$(1) L(f^1, f^2) \equiv_{df} \forall_q \left((f_q^1 = f_q^2) \vee \left(\exists_u \left((u \leq q) \& (f_u^1 \leq u f_u^2) \& (\sim (f_u^1 = f_u^2)) \right) \right) \right)$$

It is easy to check that (1) is in fact the definition of the lexicographic product.

Let $L \in \Phi$. The A -relation corresponding to L will be denoted by λ_B . Hence $(f^a)_{a \in A} \in \lambda_B$ iff $L(f^a)_{a \in A}$ holds. Let $m = (m_q)_{q \in Q} \in B$ and let λ_B satisfy the following condition (C_m) :

(C_m) : If $f^a \in B(a \in A)$, $s \in Q$, and $f^a_q = m_q$ for all $a \in A$ and all $q \in Q \setminus \{s\}$, then

$$(f^a)_{a \in A} \in \lambda_B \text{ iff } (f^a_s)_{a \in A} \in \alpha_s.$$

We shall call the set B with the A -relation λ_B satisfying (C_m) the L_m -product of (B_q, α_q) over (Q, γ) , in symbols $\prod_{q \in Q} B_q$ (shortly L -product only). It is a simple matter to check that L defined by (1) satisfies (C_m) for any $m \in B$; hence the lexicographic product is an L_m -product, in particular the cardinal product is also an L_m -product. Thus the L_m -product is a generalization of the lexicographic product.

2.

2.1. Definition. Let K, K', S , and $U_{kk'}$, ($k \in K, k' \in K'$) be sets. Let h be a mapping of $\prod_{(k,k') \in K \times K'} U_{kk'}$ into S . We define $h^*: \prod_{k \in K, k' \in K'} (\prod U_{kk'}) \rightarrow S$ by

$$h^*((u_{kk'})_{k' \in K'})_{k \in K} = h(u_{kk'})_{(k,k') \in K \times K'}$$

for any $u_{kk'} \in U_{kk'}$ ($k \in K, k' \in K'$). Obviously $h \rightarrow h^*$ is an one-to-one correspondence between the set of all mappings of $\prod_{(k,k') \in K \times K'} U_{kk'}$ into S and the set of all mappings of $\prod_{k \in K, k' \in K'} (\prod U_{kk'})$ into S . In the sequel we shall denote both mappings h and h^* by the same symbol.

2.2. Definition. Throughout this note S will be a nonempty set, σ an A -relation on S , (U, λ_U) will be a fixed L_m -product $\prod_{k \in K} U_k$ of the sets (U_k, ϱ_k) over (K, κ) and $(U', \lambda_{U'})$ will be a fixed $L_{m'}$ -product $\prod_{k' \in K'} U'_{k'}$ of the sets $(U'_{k'}, \varrho'_{k'})$ over (K', κ') (where ϱ_k and $\varrho'_{k'}$ are A -relations on U_k and $U'_{k'}$ respectively and κ and κ' are C -relations on K and K' , respectively) In the sequel we shall assume that f is an isomorphism of (U, λ_U) onto (S, σ) (that is f is a bijection such that $(g^a)_{a \in A} \in \lambda_U$ iff $(fg^a)_{a \in A} \in \sigma$), f' is an isomorphism of $(U', \lambda_{U'})$ onto (S, σ) and $fm = f'm' = n$.

If this holds then $(S, (U_k)_{k \in K}, f, n)$ and $(S, (U'_{k'})_{k' \in K'}, f', n)$ are admissible quadruples in the sense of [3] 1.1 and the mappings g_k and $g'_{k'}$ can be defined as in [3] 1.2. The following conditions are equivalent (see [3] 4 and [4] th. 1):

(α): $g_k g'_{k'} = g'_{k'} g_k$ for all $k \in K$ and $k' \in K'$,

(δ) There exist sets $U_{kk'}$ and bijections $f'_k: \prod_{k' \in K'} U_{kk'} \rightarrow U_k$ and bijections $f_k: \prod_{k \in K} U_{kk'} \rightarrow U'_k$ for every $k \in K, k' \in K'$ such that $f(f'_k)_{k \in K} = f'(f_k)_{k' \in K'}$,

where $f(f'_k)_{k \in K}$ and $f'(f_k)_{k' \in K'}$ are defined in [3] 3.5.

We shall prove now that if (α) holds, then there exist \mathcal{A} -relations $\xi_{kk'}$ on $U_{kk'}$ such that the bijections $f(f'_k)_{k \in K}$ and $f'(f_k)_{k' \in K'}$ are isomorphisms of the L-product $X = \prod_{k \in K} (\prod_{k' \in K'} U_{kk'})$ and the L-product $Y = \prod_{k' \in K'} (\prod_{k \in K} U_{kk'})$ onto (S, σ) , respectively. The proof is based on several lemmas. Throughout k and k' are elements of K and K' , respectively. We shall assume that (α) holds. We define $U_{kk'}$ as in [3] 3.8 (proof part 5) by

$$(2) \quad U_{kk'} = q_k q'_k S = q'_k q_k S (= q_k S \cap q'_k S).$$

Further let

$$(3) \quad X_k = \prod_{k' \in K'} U_{kk'}, \quad Y_{k'} = \prod_{k \in K} U_{kk'}.$$

From [3] 1.7 and 1.5(i) it follows:

2.3. Lemma. *If (α) holds and $u_{kk'} \in U_{kk'}$, then*

$$(4) \quad q_k u_{kk'} = u_{kk'}, \quad q_j u_{kk'} = n(j \in K, j \neq k),$$

$$(5) \quad q'_k u_{kk'} = u_{kk'}, \quad q'_l u_{kk'} = n(l \in K', l' \neq k').$$

2.4. Lemma. *If $t \in q_l S$, then for every $k \in K, k \neq l$, we have*

$$(6) \quad p_k f^{-1} t = p_k f^{-1} n = n_k.$$

Proof: By [3] 1.2 and 1.5 (i) $t = q_l t$, hence $f^{-1} t = f^{-1} q_l t = f^{-1} f o_l p_l f^{-1} t = o_l p_l f^{-1} t$ and the lemma follows.

From [3] 3.8 (proof part 5), 3.6 and 1.2 it follows that f'_k is given by

$$(7) \quad f'_k(u_{kk'})_{k' \in K'} = p_k f^{-1} f'(p'_k f'^{-1} u_{kk'})_{k' \in K'}.$$

Thus,

$$(8) \quad f(f'_k)_{k \in K} (u_{kk'})_{(k, k') \in K \times K'} = f(p_k f^{-1} (f'(p'_k f'^{-1} u_{kk'})_{k' \in K'}))_{k \in K}.$$

Further from [3] 3.2 and 1.2 we see that

$$(9) \quad f(f'_k)_{k \in K} = g o g'.$$

In that what follows $u_{kk'}^a$, will be elements of $U_{kk'} (a \in A)$. We put

$$(10) \quad t_k^a = (u_{kk'}^a)_{k' \in K'},$$

$$(11) \quad v_k^a = (p'_k f'^{-1} u_{kk'}^a)_{k' \in K'}.$$

Thus $t_k^a \in X_k$, $v_k^a \in U'$, and

$$(12) \quad f(f_k)_{k \in K}(u_{kk'}^a)_{(k, k') \in K \times K'} = f(p_k f^{-1} f' v_k^a)_{k \in K}.$$

2.5. Lemma. $f' v_k^a \in q_k S$, i.e.

$$(13) \quad q_k f' v_k^a = f' v_k^a.$$

Proof: As $u_{kk'}^a \in U_{kk'} = q_k S \cap q_{k'}' S$, it follows that $gog'(u_{kk'}^a)_{(k, k') \in K \times K'}$ is defined ([3] 3.3, proof part 1). From [3] 3.3, proof part 1 we obtain also that $g'(u_{kk'}^a)_{k' \in K'} \in q_k S$. But $g'(u_{kk'}^a)_{k' \in K'} = f' v_k^a$ ((11) and [3] 1.2) and (13) now follows by [3] 1.5 (i).

We note that $U_{kk'} \subseteq S$ and therefore the A -relations $\xi_{kk'}$ on $U_{kk'}$ can be defined as the restrictions of σ to $U_{kk'}$ (i.e. $\xi_{kk'} = \{g \in \sigma \mid g(A) \in U_{kk'}\}$).

2.6. Lemma. *The following conditions are equivalent:*

- (i) $(u_{kk'}^a)_{a \in A} \in \xi_{kk'}$,
- (ii) $(p_k f^{-1} u_{kk'}^a)_{a \in A} \in Q_k$,
- (iii) $(p_{k'} f'^{-1} u_{kk'}^a)_{a \in A} \in Q_{k'}$.

Proof: It follows from 2.3, (C_m) , $(C_{m'})$, $\xi_{kk'} \subseteq \sigma$, and from the fact that f and f' are isomorphisms.

2.7. Lemma. *If R is one of the atomic predicates and $k \in K$, then*

$$(14) \quad R(t_k^a)_{a \in A} \Leftrightarrow R(v_k^a)_{a \in A}.$$

Proof: By the definition in the section 1 we have to consider the following three cases:

1) Let $R = E_{a' a'' l'}(a', a'' \in A, l' \in K')$. Assume that $R(v_k^a)_{a \in A}$ holds. By (11) $R(v_k^a)_{a \in A}$ means $p_i f'^{-1} u_{ki'}^a = p_i f'^{-1} u_{ki'}^a$. In view of (5) we have $u_{ki'}^a = q_i' u_{ki'}^a = f' o_i' p_i f'^{-1} u_{ki'}^a = f' o_i' p_i f'^{-1} u_{ki'}^a = q_i' u_{ki'}^a = u_{ki'}^a$; thus by (10) $R(t_k^a)_{a \in A}$ holds. Conversely $u_{ki'}^a = u_{ki'}^a$ obviously implies $p_i f'^{-1} u_{ki'}^a = p_i f'^{-1} u_{ki'}^a$ and (14) holds.

2) Let $R = P v_l'(l' \in K', v: A \rightarrow A)$. Then $R(t_k^a)_{a \in A}$ means $(u_{ki'}^a)_{a \in A} \in \xi_{k l'}$ and $R(v_k^a)_{a \in A}$ means $(p_i f'^{-1} u_{ki'}^a)_{a \in A} \in Q_{l'}$ and the assertion follows from 2.6.

3) Let $R = W_\mu(\mu: C \rightarrow K')$. Since $t_k^a \in X_k$ and $v_k^a \in U'$ and both X_k and U' are products over the same set (K', κ') , both sides of (14) mean simply $(\mu c)_{c \in C} \in \kappa'$ and (14) is trivially satisfied.

2.8. Lemma. *If $k \in K$, then*

$$(15) \quad (t_k^a)_{a \in A} \in \lambda_{X_k} \Leftrightarrow (v_k^a)_{a \in A} \in \lambda_{U'}.$$

Proof: L is a formula constructed from the atomic predicates. Since \vee , $\&$, \sim , \exists , and \forall preserve the equivalence \Leftrightarrow , it follows from (14), that $L(t_k^a)_{a \in A}$ holds iff $L(v_k^a)_{a \in A}$ holds, hence (15) holds.

2.9. Lemma. *If R is one of the atomic predicates then*

$$(16) \quad R((t_k^a)_{k \in K})_{a \in A} \Leftrightarrow R((p_k f^{-1} f' v_k^a)_{k \in K})_{a \in A}.$$

Proof: According to the definition of R we have to consider the following three cases:

1. Let $R = E_{a' a'' l} (a', a'' \in A, l \in K)$. Then the left and right side of (16) mean $t_l^{a'} = t_l^{a''}$ and $p_l f^{-1} f' v_l^{a'} = p_l f^{-1} f' v_l^{a''}$, respectively. Since f^{-1} and f' are bijections and t_l^a determines completely v_l^a , we have $t_l^{a'} = t_l^{a''} \Rightarrow p_l f^{-1} f' v_l^{a'} = p_l f^{-1} f' v_l^{a''}$. Conversely let $p_l f^{-1} f' v_l^{a'} = p_l f^{-1} f' v_l^{a''}$. Then by 2.5 and by the definition of q_l we have $f' v_l^{a'} = q_l f' v_l^{a''} = f o_l p_l f^{-1} f' v_l^{a'} = f o_l p_l f^{-1} f' v_l^{a''} = q_l f' v_l^{a''} = f' v_l^{a''}$. But f' is a bijection and therefore $v_l^{a'} = v_l^{a''}$. By (11) we have $p_k f^{-1} u_{lk}^{a'} = p_k f^{-1} u_{lk}^{a''}$ for any $k' \in K'$. Therefore by (5) and the definition of q_k for each $k' \in K'$ we obtain $u_{lk}^{a'} = q_k u_{lk}^{a''} = f' o_k p_k f^{-1} u_{lk}^{a'} = f' o_k p_k f^{-1} u_{lk}^{a''} = q_k u_{lk}^{a''} = u_{lk}^{a''}$, i.e. $t_l^{a'} = t_l^{a''}$.

2. Let $R = P_{\nu l} (l \in K, \nu : A \rightarrow A)$. Then the left and right side of (16) mean $(t_l^a)_{a \in A} \in \lambda_{Xl}$ and $(p_l f^{-1} f' v_l^a)_{a \in A} \in \varrho_l$. According to 2.8 and in view of the fact that f' and f^{-1} are isomorphisms we have: $(t_l^a)_{a \in A} \in \lambda_{Xl} \Leftrightarrow (v_l^a)_{a \in A} \in \lambda_{U'}$ $\Leftrightarrow (f^{-1} f' v_l^a)_{a \in A} \in \lambda_U$. By 2.5, 2.4, and (C_m) we have $(f^{-1} f' v_l^a)_{a \in A} \in \lambda_U \Leftrightarrow (p_l f^{-1} f' v_l^a)_{a \in A} \in \varrho_l$ and (16) holds.

3. Let $R = W_{\mu} (\mu : C \rightarrow K)$. Since both X and U are products over (K, κ) , both sides of (16) mean simply $(\mu c)_{c \in C} \in \kappa$ and (16) is trivially satisfied.

2.10. Lemma. $f(f_k)_{k \in K}$ is an isomorphism of X onto S .

Proof: $f(f_k)$ is a bijection of X onto S . f is an isomorphism, hence in view of (8), (10), and (11) we have to prove only that

$$(17) \quad L((t_k^a)_{k \in K})_{a \in A} \Leftrightarrow L((p_k f^{-1} f' v_k^a)_{k \in K})_{a \in A}.$$

L is a formula constructed from atomic predicates. Since \vee , $\&$, \sim , \exists , and \forall preserve the equivalence \Leftrightarrow , (17) is a consequence of (16).

By symmetry we have a similar statement for $f'(f_k)_{k' \in K'}$. Thus we have

2.11. Theorem. Let S be a set, σ an A -relation on S , (U, λ_U) and $(U', \lambda_{U'})$ be an L_m -product of the sets (U_k, ϱ_k) over (K, κ) and an L_m -product of the sets $(U_{k'}, \varrho_{k'})$ over (K', κ') respectively (where ϱ_k and $\varrho_{k'}$ are A -relations on U_k and $U_{k'}$ respectively and κ and κ' are C -relations on K and K' respectively). Further let f and f' be isomorphisms of (U, λ_U) and $(U', \lambda_{U'})$ onto (S, σ) respectively such that $f m = f' m' = n$. Then the following assertions are equivalent:

(α) For every $k \in K$ and $k' \in K'$ the mappings q_k and $q_{k'}$ (determined by n) satisfy

$$q_k q_{k'} = q_k q_{k'}.$$

δ^*) For every $k \in K$ and $k' \in K'$ there exists a set $U_{kk'}$ and an A -relation $\xi_{kk'}$ on $U_{kk'}$; for every $k \in K$ there exists a bijection $f_k : \prod_{k' \in K'} U_{kk'} \rightarrow U_k$, and for every $k' \in K'$ there exists a bijection $f_{k'} : \prod_{k \in K} U_{kk'} \rightarrow U_{k'}$ such that $f(f_k)_{k \in K}$ and $f'(f_{k'})_{k' \in K'}$ are isomorphisms of the L -product $\prod_{k \in K} (\prod_{k' \in K'} U_{kk'})$ and of the L -product $\prod_{k' \in K'} (\prod_{k \in K} U_{kk'})$ onto (S, σ) respectively.

3.

Let S be a set. Further let Ω be a set and let $\mathcal{A} = \{A_\omega \mid \omega \in \Omega\}$ be a system of nonempty sets. If σ^ω are any A_ω -relations on S , then the system $\{\sigma^\omega \mid \omega \in \Omega\}$ will be said to be an \mathcal{A} -relational structure on S . Let $\{\rho_k^\omega \mid \omega \in \Omega\}$ be an \mathcal{A} -relational structure on U_k for each $k \in K$. If the L -product of (U_k, ρ_k^ω) over (K, \varkappa) is denoted by (U, λ_V^ω) for each $\omega \in \Omega$, then $\{\lambda_V^\omega \mid \omega \in \Omega\}$ is obviously an \mathcal{A} -relational structure on U . We say that f is an isomorphism of the \mathcal{A} -relational structure $\{\lambda_V^\omega \mid \omega \in \Omega\}$ onto the \mathcal{A} -relational structure $\{\sigma^\omega \mid \omega \in \Omega\}$ iff f is an isomorphism of the relation λ_V^ω onto the relation σ^ω for each $\omega \in \Omega$.

It is easy to see that 2.11 remains valid if we replace the relations σ, ρ_k, ρ_k' by \mathcal{A} -relational structures.

Let ρ be an A -relation on M and let $a_0 \in A$. We say that ρ is an operation on M iff for any $(f_a)_{a \in A} \in \rho$ and $(g_a)_{a \in A} \in \rho$ we have: $f_a = g_a$ for every $a \in A, a \neq a_0$ implies $f_{a_0} = g_{a_0}$. Hence a finitary n -ary operation is a special case of $(n + 1)$ -ary relation. Our definition includes also partial and infinitary operations. From this it follows that universal algebras may be regarded as a special case of relational structures.

In the sequel we shall restrict ourselves to the use of (full) direct product of algebras. A subalgebra with a single element $\{n\}$ is termed a trivial subalgebra (for reference see e.g. [1]). Obviously an isomorphism carries a trivial subalgebra onto a trivial subalgebra and (C_n) holds. Hence we have (see also [4] Theorem 2):

3.1. Theorem. Let Ω be an operator domain. Let S be an Ω -algebra with the trivial subalgebra $\{n\}$ and let $U_k (k \in K)$ and $U_{k'} (k' \in K')$ be Ω -algebras. Further let f and f' be isomorphisms of the direct products $\prod_{k \in K} U_k$ and $\prod_{k' \in K'} U_{k'}$ onto S respectively. Then the following assertions are equivalent:

(α) For every $k \in K$ and $k' \in K'$ the mappings q_k and $q_{k'}$ (determined by n) satisfy

$$q_k q_{k'} = q_{k'} q_k.$$

(δ') For every $k \in K$ and $k' \in K'$ there exists an Ω -algebra $U_{kk'}$; for every $k \in K$ there exists an isomorphism $f'_k : \prod_{k' \in K'} U_{kk'} \rightarrow U_k$ and for every

$k' \in K'$ there exists an isomorphism $f_{k'} : \prod_{k \in K} U_{kk'} \rightarrow U_{k'}$ such that $f(f_k)_{k \in K}$ and $f'(f_{k'})_{k' \in K'}$ are isomorphisms of $\prod_{(k,k') \in K \times K'} U_{kk'}$ onto S respectively.

3.2. Remark. Hashimoto has proved in [2] that if σ is a binary, reflexive, antisymmetric and connected relation on S , then the condition (α) is satisfied for cardinal products. This is not true for lexicographic products as the following very simple example shows.

Let $h > 0$ be an integer. By h we understand the chain $0 < 1 < \dots < h - 1$. Let $K = K' = 2$, $U_0 = U'_1 = 2$, $U_1 = U'_0 = 3$. Then $U = \prod_{k \in K} U_k$ and $U' = \prod_{k' \in K'} U_{k'}$ are chains with 6 elements and hence both are isomorphic to $S = 6$. Let $n = 0$, $s = 3$ ($\in S$). Then $q_0 3 = f_{00} p_0(1, 0) = f(1, 0) = 3$ and similarly $q'_1 3 = f'_{0'_1} p'_1(1, 1) = f'(0, 1) = 1$. Thus $q'_1 q_0 3 = q'_1 3 = 1$ and $q_0 q'_1 3 = q_0 1 = f_{00} p_0(0, 1) = f(0, 0) = 0$, i.e. $q_0 q'_1 3 \neq q'_1 q_0 3$.

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