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CONCRETE QUANTUM LOGICS WITH GENERALISED  
COMPATIBILITY

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*Abstract.* We present three results stating when a concrete (= set-representable) quantum logic with covering properties (generalization of compatibility) has to be a Boolean algebra. These results complete and generalize some previous results [3, 5] and answer partially a question posed in [2].

*Keywords:* Boolean algebra, concrete quantum logic, covering, Jauch-Piron state, orthocompleteness

*MSC 1991:* 03G12, 81P10

## 1. BASIC NOTIONS

Let us recall the main notion we shall deal with in this paper.

**Definition 1.1.** A *concrete logic* is a pair  $(X, L)$ , where  $X \neq \emptyset$  and  $L \subset \exp X$  such that

- (1)  $\emptyset \in L$ ;
- (2)  $A^c = X \setminus A \in L$  whenever  $A \in L$ ;
- (3)  $\bigcup M \in L$  whenever  $M \subset L$  is a finite set of mutually disjoint elements.

A *concrete  $\sigma$ -logic* is a concrete logic  $(X, L)$  such that

- (3 $\sigma$ )  $\bigcup M \in L$  whenever  $M \subset L$  is a countable set of mutually disjoint elements.

Let us note that the above definition is not given in the most efficient way. Indeed, since  $\emptyset$  is a finite set of mutually disjoint elements and  $\bigcup \emptyset = \emptyset$ , condition (1) follows from condition (3). Moreover, it is obvious that condition (3) follows from condition (3 $\sigma$ ).

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The following lemma will be useful in the sequel. First, let us observe that if  $A, B \in L$  and  $A \subset B$ , then  $B \setminus A = (A \cup B^c)^c \in L$  for every concrete logic  $(X, L)$ .

**Lemma 1.2.** *Let  $(X, L)$  be a concrete  $\sigma$ -logic and let  $A_i \in L$  ( $i = 1, 2, \dots$ ) be such that  $A_1 \supset A_2 \supset \dots$ . Then  $\bigcap_{i=1}^{\infty} A_i \in L$ .*

**Proof.** The elements  $A_i \setminus A_{i+1} \in L$  ( $i = 1, 2, \dots$ ) are mutually orthogonal, hence  $\bigcup_{i=1}^{\infty} (A_i \setminus A_{i+1}) \in L$  and  $\bigcap_{i=1}^{\infty} A_i = A_1 \setminus \bigcup_{i=1}^{\infty} (A_i \setminus A_{i+1}) \in L$ .  $\square$

## 2. COVERING PROPERTIES

**Definition 2.1.** Let  $(X, L)$  be a concrete logic,  $Y \subset X$  and let  $n$  be a natural number. A *covering* of  $Y$  is a set  $M \subset L$  such that  $Y = \bigcup M$ . A covering  $M$  is an *n-covering* if  $\text{card } M \leq n$ .

We say that  $(X, L)$  has the *n-covering property* (*finite covering property*, resp.) if for every  $A, B \in L$  there is an *n-covering* (finite covering, resp.) of  $A \cap B$ .

It is well-known that a concrete logic  $(X, L)$  is a Boolean algebra if and only if  $A \cap B \in L$  for every  $A, B \in L$ , i. e. if and only if  $(X, L)$  has the 1-covering property. Thus, the notions of *n-covering property* (finite covering property), introduced in [3], are generalizations of compatibility in Boolean algebras.

The next lemma will be used in the sequel.

**Lemma 2.2.** *Let  $(X, L)$  be a concrete logic with the finite covering property. Then for every finite set  $F \subset L$  there is a finite covering  $G \subset L$  of  $\bigcap F$ .*

**Proof.** Let us proceed by induction. First, if  $F$  is a one-element subset of  $L$  (empty set, resp.), then we can put  $G = F$  ( $G = \{X\}$ , resp.).

Now, let us suppose that there is a natural number  $n \geq 1$  such that the lemma holds for every  $F \subset L$  with  $\text{card } F = n$ . Let  $F \subset L$  with  $\text{card } F = n + 1$  and let  $A \in F$ . According to the previous assumption, there is a finite covering  $G \subset L$  of  $\bigcap (F \setminus \{A\})$ . According to the finite covering property, for every  $B \in G$  there is a finite covering  $G_B \subset L$  of  $A \cap B$ . Thus,  $\bigcup_{B \in G} G_B \subset L$  is a finite covering of  $\bigcap F$ .  $\square$

Before we present the main result of this section, let us prove the following technical lemma.

**Lemma 2.3.** *Let  $(X, L)$  be a concrete  $\sigma$ -logic and let  $m, n \geq 2$  be natural numbers such that  $m \leq n + 1$ . Let us suppose that for every set  $F \subset L$  with  $\text{card } F \leq n$  there is an  $m$ -coverig  $G \subset L$  of  $\bigcap F$ . Then for every set  $F \subset L$  with  $\text{card } F \leq n$  there is an  $(m - 1)$ -covering  $G \subset L$  of  $\bigcap F$ .*

*Proof.* Let  $F \subset L$  with  $\text{card } F \leq n$ . Let us define by induction sequences  $(A_{i1}, \dots, A_{in}) \in L^n$ ,  $(B_{i0}, \dots, B_{im}) \in L^{n+1}$  ( $i = 1, 2, \dots$ ) as follows: Let  $(A_{11}, \dots, A_{1n})$  be such that  $F = \{A_{11}, \dots, A_{1n}\}$ . If  $(A_{i1}, \dots, A_{in}) \in L^n$  is defined for a natural number  $i \geq 1$  then let us take  $(B_{i0}, \dots, B_{im}) \in L^{n+1}$  such that  $B_{ij} = \emptyset$  for  $j \geq m$  and  $\bigcap_{j=1}^n A_{ij} = \bigcup_{j=0}^n B_{ij}$  and let us put  $A_{i+1,j} = A_{ij} \setminus B_{ij}$  ( $j \in \{1, \dots, n\}$ ).

Let us denote

$$B_0 = \bigcap_{i=1}^{\infty} B_{i0}, \quad B_j = \bigcup_{i=1}^{\infty} B_{ij}, \quad j \in \{1, \dots, n\}.$$

It is easy to see that the elements  $B_{1j}, B_{2j}, \dots$  ( $j \in \{1, \dots, n\}$ ) are mutually disjoint, hence  $B_j \in L$  for every  $j \in \{1, \dots, n\}$ . Moreover,  $B_m = \dots = B_n = \emptyset$ . Further,

$$B_{i0} \supset \bigcap_{j=1}^n A_{i+1,j} \supset B_{i+1,0} \quad (i = 1, 2, \dots).$$

Hence, according to Lemma 1.2.  $B_0 \in L$ , too. Since

$$\bigcap F = B_0 \cup B_1 \cup \dots \cup B_{m-1}$$

and since  $B_0 \cup B_1 \in L$  ( $B_0 \cap B_1 = \emptyset$ ), the proof is complete.  $\square$

**Theorem 2.4.** *Let  $(X, L)$  be a concrete  $\sigma$ -logic. Let us suppose that there is a natural number  $n \geq 2$  such that for any set  $F \subset L$  with  $\text{card } F \leq n$  there is an  $(n + 1)$ -covering of  $\bigcap F$ . Then  $(X, L)$  is a Boolean algebra.*

*Proof.* Using Lemma 2.3  $n$ -times, we obtain that  $(X, L)$  has the 1-covering property, i. e.  $(X, L)$  is a Boolean algebra.  $\square$

**Corollary 2.5.** *Every concrete  $\sigma$ -logic with the 3-covering property is a Boolean algebra.*

This corollary generalizes [3, Proposition 4.6], where an analogous result is stated for concrete  $\sigma$ -logics with the 2-covering property.

### 3. COVERING PROPERTIES AND JAUCH-PIRON STATES

**Definition 3.1.** Let  $(X, L)$  be a concrete logic. A *state* on  $(X, L)$  is a mapping  $s: L \rightarrow [0, 1]$  such that

- (1)  $s(X) = 1$ ;
- (2)  $s(\bigcup M) = \sum_{A \in M} s(A)$  whenever  $M \subset L$  is a finite set of mutually disjoint elements.

A state  $s$  on  $(X, L)$  is called *Jauch-Piron* if for every  $A, B \in L$  with  $s(A) = s(B) = 1$  there is a  $C \in L$  such that  $C \subset A \cap B$  and  $s(C) = 1$ .

It is easy to see that  $s(\emptyset) = 0$  and  $s(A^c) = 1 - s(A)$  for every state  $s$  on a concrete logic  $(X, L)$  and for every  $A \in L \setminus \{\emptyset\}$ . Further, for every concrete logic  $(X, L)$ , every point  $x \in X$  carries a two-valued state  $s_x$  on  $(X, L)$  defined by

$$s_x(A) = \begin{cases} 1, & \text{if } x \in A, \\ 0, & \text{if } x \notin A. \end{cases}$$

Before we present the main result of this section, we need the following definition.

**Definition 3.2.** Let  $(X, L)$  be a concrete logic and let  $M, N \subset L$  be two coverings of  $Y \subset X$ . We say that  $N$  is a *coarsing* of  $M$  if for every  $A \in M$  there is a  $B \in N$  such that  $A \subset B$ .

**Theorem 3.3.** Let  $(X, L)$  be a concrete logic such that every state on it is Jauch-Piron. Let us suppose that for every  $A, B \in L$  every covering of  $A \cap B$  admits a countable coarsing. Then  $L$  is a Boolean algebra.

*Proof.* It suffices to prove that  $A \cap B \in L$  for every  $A, B \in L$ . Let  $A, B \in L$ . If  $A \cap B = \emptyset$ , the proof is complete. Let us suppose that  $A \cap B \neq \emptyset$ . Then  $S_{A,B} = \{s; s \text{ is a state on } (X, L) \text{ with } s(A) = s(B) = 1\}$  is nonempty (every point  $x \in A \cap B$  carries a two-valued state  $s_x \in S_{A,B}$ ). Since every state on  $(X, L)$  is Jauch-Piron, for every  $s \in S_{A,B}$  there is a  $C_s \in L$  such that  $s(C_s) = 1$ . Let us take a countable coarsing  $M$  of the covering  $\{C_s; s \in S_{A,B}\}$  of  $A \cap B$ , a countable set  $Y \subset A \cap B$  such that  $Y \cap (C \setminus D) \neq \emptyset$  for every  $C, D \in M$  with  $C \setminus D \neq \emptyset$  and, finally, a state  $s$  that is a  $\sigma$ -convex combination (with non-zero coefficients) of all  $s_y$  ( $y \in Y$ ). Since  $s \in S_{A,B}$ , there is a  $D_s \in M$  such that  $s(D_s) = 1$ . Thus,  $D_s \supset Y$  and therefore  $A \cap B = \bigcup M = D_s \in L$ .  $\square$

Theorem 3.3 seems to be independent of the previous results in [3, 4, 7], nevertheless it has corollaries that were obtained using quite a different techniques. (Let us note that a unifying look at these attempts is presented in [8].) The following corollary was obtained (in a more general form) in [4].

**Corollary 3.4.** *Every countable concrete logic such that every state on it is Jauch-Piron is a Boolean algebra.*

The next corollary of Theorem 3.3 was obtained (in a more general form) in [7].

**Corollary 3.5.** *Let  $(X, L)$  be a concrete logic such that every state on it is Jauch-Piron. Let us suppose that  $(X, L)$  contains only countably many maximal Boolean subalgebras and these are complete. Then  $(X, L)$  is a Boolean algebra.*

**PROOF.** It is easy to see that for every  $A, B \in L$  every covering of  $A \cap B$  admits a countable coarsening.  $\square$

#### 4. COVERING PROPERTIES AND ORTHOCOMPLETENESS

**Definition 4.1.** Let  $\alpha$  be a cardinal number. A concrete logic  $(X, L)$  is called  $\alpha$ -orthocomplete if  $\bigvee M \in L$  (supremum with respect to inclusion) whenever  $M \subset L$  is a set of mutually disjoint elements with  $\text{card } M \leq \alpha$ .

It is obvious that condition  $(3\sigma)$  of Definition 1.1 implies that a concrete  $\sigma$ -logic is  $\omega_0$ -orthocomplete ( $\omega_0$  denotes the countable cardinal)—this is usually denoted as  $\sigma$ -orthocompleteness.

The following theorem generalizes a result from [5] and answers partially a question posed in [2].

**Theorem 4.2.** *Every  $c$ -orthocomplete ( $c$  denotes the cardinality of real numbers) concrete  $\sigma$ -logic with the finite covering property is a Boolean algebra.*

**PROOF.** Let  $(X, L)$  be a concrete  $\sigma$ -logic with the finite covering property and let  $A, B \in L$ . It suffices to prove that  $A \cap B \in L$ . Let us define by induction finite subsets  $F_i$  ( $i = 1, 2, \dots$ ) of  $L$  as follows: First,  $F_1 \subset L$  is a finite covering of  $A \cap B$ . Now, let a finite set  $F_i = \{A_1, \dots, A_n\} \subset L$  be defined for a natural number  $i \geq 1$ . Let us denote by  $G_i$  the set of all intersections of the form  $A_1^{e_1} \cap \dots \cap A_n^{e_n}$ , where  $(e_1, \dots, e_n) \in \{-1, 1\}^n \setminus \{-1\}^n$  and  $A_j^1 = A_j, A_j^{-1} = X \setminus A_j$  ( $j = 1, \dots, n$ ).  $G_i$  is a finite set of mutually disjoint subsets of  $X$  such that  $\bigcap F_i = \bigcup G_i$ . According to Lemma 2.2, for every  $Y \in G_i$  there is a finite covering  $G_Y \subset L$  of  $Y$ . Let us put  $F_{i+1} = \bigcup_{Y \in G_i} G_Y$ .

Let us consider all sequences  $C_1, C_2, \dots$  such that  $C_i \in F_i$  ( $i = 1, 2, \dots$ ) and  $C_1 \supset C_2 \supset \dots$ . According to Lemma 1.2,  $\bigcap_{i=1}^{\infty} C_i \in L$  for each such sequence. Hence, we have obtained at most the continuum of mutually disjoint elements of  $L$  such that their union is  $A \cap B$ . Since their supremum exists, it is equal to  $A \cap B$ . Thus,  $A \cap B \in L$ .  $\square$

Before we present a corollary of Theorem 4.2, let us recall a result connecting the covering properties with Jauch-Piron states [3, Theorem 3.5].

**Theorem 4.3.** *Let  $(X, L)$  be a concrete logic such that every two-valued state on it is Jauch-Piron. Then  $(X, L)$  has the finite covering property.*

**Corollary 4.4.** *Every  $c$ -orthocomplete concrete  $\sigma$ -logic such that every two-valued state on it is Jauch-Piron is a Boolean algebra.*

**Proof.** It follows from Theorem 4.3 and Theorem 4.2. □

**Remark 4.5.** The above corollary can be stated in the following (more general) way: Every  $c$ -orthocomplete quantum  $\sigma$ -logic with a closed full set of two-valued Jauch-Piron  $\sigma$ -states is a Boolean algebra. Indeed, concrete  $\sigma$ -logics are exactly representations of quantum  $\sigma$ -logics with a full set of two-valued  $\sigma$ -states (see e.g. [1, 6]) and Theorem 4.3 can be stated for quantum logics with a closed full set of two-valued Jauch-Piron states (the set of two-valued states is closed in the product topology on  $[0, 1]^L$ ).

The following question (posed in [2]) remains open. Here we have given the negative answer in the case that the concrete logic in question is also  $c$ -orthocomplete.

**Question 4.6.** Is there a concrete  $\sigma$ -logic that is not a Boolean algebra such that every state on it is Jauch-Piron?

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