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ON A GENERALIZATION OF THE MATROID

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The notation of the semimatroid is known from literature (see [2]). A semimatroid is a pair $H = \langle X, \mathcal{B} \rangle$, where X is a non empty finite set and \mathcal{B} is a non empty anti-hereditary family of subsets of X , which means that \mathcal{B} satisfies the condition

$$(i) \quad B \in \mathcal{B} \wedge A \subseteq B \Rightarrow A \notin \mathcal{B}.$$

The sets from \mathcal{B} will be called *bases of the semimatroid* H . A semimatroid $H = \langle X, \mathcal{B} \rangle$ satisfying the condition

$$(e) \quad B_1, B_2 \in \mathcal{B} \Rightarrow \bigwedge_{x \in B_1 \setminus B_2} \bigvee_{y \in B_2 \setminus B_1} [(B_1 \setminus \{x\}) \cup \{y\} \in \mathcal{B}]$$

is called a *matroid* (see [3]).

In this paper we consider a generalization of the condition (e), namely: a semimatroid $H = \langle X, \mathcal{B} \rangle$ will be called an *e*-semimatroid* if H satisfies the condition

$$(e^*) \quad B_1, B_2 \in \mathcal{B} \Rightarrow \bigwedge_{x_1 \in B_1 \setminus B_2} \bigvee_{y_1 \in B_1} \bigvee_{x_2, y_2 \in B_2} [(B_1 \setminus \{x_1, y_1\}) \cup \{x_2, y_2\} \in \mathcal{B} \wedge (x_1 = y_1 \Rightarrow x_2 = y_2)].$$

It is easy to verify that any matroid is an e*-semimatroid. On the other hand, there are e*-semimatroids that are not matroids, which is shown by the following example.

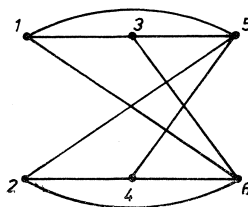


Figure 1.

Example 1. Let us take the simple graph G (see [1]) in Fig. 1. Let us consider a semimatroid $H = \langle X, \mathcal{B} \rangle$, where $X = \{1, 2, 3, 4, 5, 6\}$ and \mathcal{B} is

the set of all cliques of G , i.e. the sets of vertices of maximal complete subgraphs of G . Therefore $\mathcal{B} = \{\{1, 3, 5\}, \{1, 3, 6\}, \{2, 4, 5\}, \{2, 4, 6\}\}$. Observe that H satisfies (e*) but not (e). We see that some semimatroids generated by graphs can be e*-semimatroids which was the reason for the authors to consider e*-semimatroids.

In this paper we prove in Section 1 that any two bases of an e*-semimatroid have the same number of elements (Theorem 1).

In Section 2 (Theorem 2) we show the following result: *Let $G = (U, X)$ be a simple connected graph. Let T be the set of edges of a spanning tree of G , let T^* be obtained from T by removing one pendant edge. Denote by \mathcal{T}^* the family of all sets of the form T^* . Then the pair $H = \langle X, \mathcal{T}^* \rangle$ is an e*-semimatroid but not necessarily a matroid.*

In Section 3 we give a representation of graphs in which cliques form a matroid and produce examples of graphs in which cliques form an e*-semimatroid.

1

Let $H = \langle X, \mathcal{B} \rangle$ be an e*-semimatroid. We shall consider the following condition:

$$(*) \quad B_1, B_2 \in \mathcal{B}, \quad B_3 = (B_1 \setminus \{x_1, y_1\}) \cup \{x_2, y_2\} \in \mathcal{B},$$

where $x_1 \in B_1 \setminus B_2$, $y_1 \in B_1$, $x_2, y_2 \in B_2$ and $(x_1 = y_1 \Rightarrow x_2 = y_2)$.

Lemma 1. *If (*) holds then $\{x_2, y_2\} \not\subseteq B_1 \cap B_2$.*

Proof. Let $x_2, y_2 \in B_1 \cap B_2$, then $B_1 \setminus \{x_1, y_1\} = B_3 \not\subseteq B_1$. Hence and by (i) $B_3 \notin \mathcal{B}$ – a contradiction.

Lemma 2. *Let (*) hold, let $B_3 \setminus B_1 = \{y_2\}$ and $x_1 \neq y_1$, then $x_2 = y_1$.*

Proof. Suppose $x_2 \neq y_1$. Then $y_1 \notin B_3$. We apply (e*) to the bases B_3 and B_1 and to the element y_2 . So there exist $z_2 \in B_3$ and $u_1, v_1 \in B_1$ such that $(B_3 \setminus \{y_2, z_2\}) \cup \{u_1, v_1\} = B_4 \in \mathcal{B}$ and $(y_2 = z_2 \Rightarrow u_1 = v_1)$. We shall show that $B_4 \not\subseteq B_1$ thus obtaining a contradiction $B_4 \in \mathcal{B}$. Let $z_2 \in B_1$. Since $x_1, y_1 \notin B_3$ and $B_3 \setminus B_1 = \{y_2\}$ so $B_4 = (B_1 \setminus \{x_1, y_1, z_2\}) \cup \{u_1, v_1\}$. Observe that $z_2 \neq x_1$ and $z_2 \neq y_1$, since $x_1, y_1 \notin B_3$ and $z_2 \in B_3$. Moreover, $x_1 \neq y_1$, hence $B_4 \not\subseteq B_1$. Let now $z_2 \notin B_1$. Then $z_2 \in B_3 \setminus B_1$, so $z_2 = y_2$. Thus we have $u_1 = v_1$, hence $B_4 = (B_1 \setminus \{x_1, y_1\}) \cup \{u_1\} \not\subseteq B_1$.

Lemma 3. *If (*) holds then we have exactly one of the following three possibilities:*

- (1) $x_1 \in B_1 \setminus B_2$, $y_1, x_2 \in B_1 \cap B_2$, $y_2 \in B_2 \setminus B_1$ and $x_2 = y_1$;
- (2) $x_1 \in B_1 \setminus B_2$, $y_1 \in B_1 \cap B_2$, $x_2, y_2 \in B_2 \setminus B_1$ and $x_2 \neq y_2$;
- (3) $x_1, y_1 \in B_1 \setminus B_2$, $x_2, y_2 \in B_2 \setminus B_1$ and $(x_1 = y_1$ and $x_2 = y_2$ or $x_1 \neq y_1$ and $x_2 \neq y_2)$.

Proof. Consider all the possibilities for the condition (e*):

- (4) $y_1 \in B_1 \cap B_2, x_2, y_2 \in B_1 \cap B_2;$
- (5) $y_1 \in B_1 \cap B_2, x_2 \in B_1 \cap B_2, y_2 \notin B_1 \cap B_2;$
- (6) $y_1 \in B_1 \cap B_2, x_2, y_2 \notin B_1 \cap B_2;$
- (7) $y_1 \notin B_1 \cap B_2, x_2, y_2 \in B_1 \cap B_2;$
- (8) $y_1 \notin B_1 \cap B_2, x_2 \in B_1 \cap B_2, y_2 \notin B_1 \cap B_2;$
- (9) $y_1 \notin B_1 \cap B_2, x_2, y_2 \notin B_1 \cap B_2.$

By Lemma 1 the cases (4), (7) cannot hold. The case (5) gives (1) by Lemma 2. The case (8) cannot hold. In fact, if $x_2 \in B_1 \cap B_2$ and $y_2 \notin B_1 \cap B_2$ then $x_2 \neq y_2$. Hence by (*) we have $B_3 \setminus B_1 = \{y_2\}$ and $x_1 \neq y_1$. In view of Lemma 2 we get $x_2 = y_1$ which is impossible. Consider the case (6). Let $y_1 \in B_1 \cap B_2, x_2, y_2 \notin B_1 \cap B_2$ and (*) hold. Then $x_2 \neq y_2$. In fact, if $x_2 = y_2$ then $B_3 \setminus B_1 = \{y_2\}$. Since $x_1 \neq y_1$ so by Lemma 2 we have $x_2 = y_1$ which is impossible. So we have the possibility (2).

Consider the case (9). Let $y_1, x_2, y_2 \notin B_1 \cap B_2$ and (*) hold. If $x_1 = y_1$ then $x_2 = y_2$ by (*). Suppose that $x_1 \neq y_1$ and $x_2 = y_2$. Then $B_3 \setminus B_1 = \{y_1\}$ and by Lemma 2 $x_2 = y_1$ which cannot hold. So we have $x_1 \neq y_1 \Rightarrow x_2 \neq y_2$. In the case (9) we have the possibility (3).

Corollary 1. *If (*) holds then $|B_3| = |B_1|$.*

Proof. It follows from Lemma 3 that the number of elements rejected from B_1 is equal to the number of elements added to B_1 .

Theorem 1. *Any two bases of an e*-semimatroid have the same number of elements.*

Proof. Let $H = \langle X, \mathcal{B} \rangle$ be an e*-semimatroid and $B_1, B_2 \in \mathcal{B}$. If $B_1 \setminus B_2 \neq \emptyset$ and $x_1 \in B_1 \setminus B_2$ then we can form the basis B_3 as in (*) and by Corollary 1 we have $|B_3| = |B_1|$. By Lemma 3 we obtain that the basis B_3 arises by deleting at least one element from the set $B_1 \setminus B_2$ and adding at least one element from the set $B_2 \setminus B_1$. If $B_3 \setminus B_2 \neq \emptyset$ then we can form the basis B_4 for the bases B_3 and B_2 analogously as we formed the basis B_3 for the bases B_1 and B_2 in (*).

Then by Corollary 1 we get $|B_4| = |B_3| = |B_1|$. Observe that $B_3 \setminus B_2 \subset B_1 \setminus B_2$ and B_4 arises by deleting at least one new element from the set $B_1 \setminus B_2$ and adding at least one new element from the set $B_2 \setminus B_1$. After a finite number of steps we get a basis B_k such that $B_k \subset B_2$ and $|B_k| = |B_{k-1}| = |B_{k-2}| = \dots = |B_3| = |B_1|$. By (i) we have $B_k = B_2$ which completes the proof.

2

Let $G = (U, X)$ be a simple connected graph. It is known that a pair $\langle X, \mathcal{T} \rangle$, where \mathcal{T} is the set of all spanning trees of G , is matroid (see [1]). L. Szamkołowicz asked which subsets of X form an e*-semimatroid. We answer this question in the following theorem.

Theorem 2. Let $G = (U, X)$ be a simple connected graph. Let T be the set of edges of a spanning tree of G , and let T^* be obtained from T by removing one pendant edge. Let \mathcal{T}^* denote the family of all sets of the form T^* . Then the pair $H = \langle X, \mathcal{T}^* \rangle$ is an e^* -semimatroid but not necessarily a matroid.

Proof. If $|U| \leq 3$ then any spanning tree has at most two edges and any T^* has at most one edge and (e^*) is satisfied. Suppose $|U| \geq 4$. Let T^* be obtained from T by removing a pendant edge. Denote the removed edge by $p(T)$. Denote by $i(T^*)$ the vertex of $p(T)$ which becomes an isolated vertex after removing $p(T)$ from the tree T . Observe that \mathcal{T}^* satisfies the condition (i) of the semimatroid. We shall show that (e^*) holds. Let $T_1^*, T_2^* \in \mathcal{T}^*$ and $x_1 \in T_1^*$. The graph $(U, T_1^* \setminus \{x_1\})$ has three components $i(T_1^*), K_1, K_2$.

If there exists $x_2 \in T_2^*$ such that $x_2 = \{v_1, v_2\}$ and $v_1 \in K_1, v_2 \in K_2$ then $(T_1^* \setminus \{x_1\}) \cup \{x_2\} = T_3^* \in \mathcal{T}^*$ since $T_3^* \cup \{p(T_1^*)\} \in \mathcal{T}$. If such an edge x_2 does not exist then necessarily $i(T_1^*) \neq i(T_2^*)$ since otherwise we have three components in the graph (U, T_2^*) . If $|K_1| = 1$ and $K_1 = \{i(T_2^*)\}$ then there exists in T_2^* an edge $x_2 = \{u, i(T_1^*)\}$ with $u \in K_2$. Hence $(T_1^* \setminus \{x_1\}) \cup \{x_2\} \in \mathcal{T}^*$. We have the analogous situation if $K_2 = \{i(T_2^*)\}$. In the remaining case there exist in T_2^* edges $\{u, i(T_1^*)\}$ and $\{i(T_1^*), v\}$ with $u \in K_1, v \in K_2$. Putting $x_2 = \{u, i(T_1^*)\}, y_2 = \{i(T_1^*), v\}$ and taking for y_1 an arbitrary pendant edge of T_1^* different from x_1 we obtain $T_3^* = (T_1^* \setminus \{x_1, y_1\}) \cup \{x_2, y_2\}$. Such an edge y_1 exists since one of the components K_1, K_2 has more than one vertex. Obviously $T_3^* \in \mathcal{T}^*$ as $(T_1^* \setminus \{x_1\}) \cup \{x_1, y_2\}$ is a spanning tree.

Consider now the graph in Fig. 2.

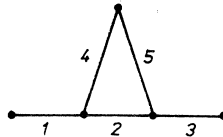


Figure 2.

Denote $T_1^* = \{1, 2, 3\}, T_2^* = \{3, 4, 5\}$. For $x_1 = 2$ there does not exist $x_2 \in T_2^* \setminus T_1^*$ such that $(T_1^* \setminus \{x_1\}) \cup \{x_2\} \in \mathcal{T}^*$. This proves the second part of Theorem 2.

3

Now we shall consider the problem when the cliques of a graph form a matroid or an e^* -semimatroid.

It is known that

(α) if $\langle X, \mathcal{B} \rangle$ is a matroid then the following condition is satisfied:

$$\bigwedge_{B_1, B_2 \in \mathcal{B}} \bigwedge_{y \in B_2 \setminus B_1} \bigvee_{x \in B_1 \setminus B_2} [(B_1 \setminus \{x\}) \cup \{y\} \in \mathcal{B}].$$

Let $G = (U, X)$ be a simple graph. We denote the set of the cliques of G by \mathcal{K}_G .

Theorem 3. The pair $H = \langle U, \mathcal{H}_G \rangle$ is a matroid iff there exists a partition $\{U_i\}_{i \in I}$ of the set U such that

$$(**) \quad \bigwedge_{u \in U_s} \bigwedge_{v \in U_t} [\{u, v\} \in X \Leftrightarrow s \neq t] \quad \text{where } s, t \in I;$$

in other words G is a complete 1-partite graph.

Proof. Sufficiency is obvious since the only cliques of G are the sets of the form B^* where $\bigwedge_{i \in I} |B^* \cap U_i| = 1$.

Proof of necessity. Let $H = \langle U, \mathcal{H}_G \rangle$ be a matroid and K_0 a fixed basis of H , i.e. a fixed clique of G . For every $a \in K_0$ we define the set $U_a = \{u \in U; (K_0 \setminus \{a\}) \cup \{u\} \in \mathcal{H}_G\}$. We shall show that the family $\{U_a\}_{a \in U}$ is a partition of U . Obviously for any $a \in K_0$ we have $U_a \neq \emptyset$ since $a \in U_a$. If $a \neq b$, $a, b \in K_0$ then $U_a \cap U_b = \emptyset$.

Suppose on the contrary that a vertex u in the graph G is connected with all vertices of the clique K_0 . This, however, contradicts the maximality of K_0 .

$$\bigwedge_{u \in U} \bigvee_{a \in K_0} u \in U_a.$$

In fact, if $u \in K_0$, then $u \in U_u$. Let $u \in U \setminus K_0$. Obviously $\{u\}$ can be extended to the maximal complete subgraph M of the graph G . So M is a clique and $M \in \mathcal{H}_G$. Applying Lemma 4 with $y = u$ and the cliques K_0 and M we find $a \in K_0$ such that $(K_0 \setminus \{a\}) \cup \{u\}$ is a clique, hence $u \in U_a$.

Now we shall show the validity of the condition (**). We shall show

$$(10) \quad \text{if } u \in U_s, v \in U_t, \{u, v\} \in X \text{ then } s \neq t.$$

Suppose $u \in U_s, v \in U_t, \{u, v\} \in X$ and $s = t$. Since $u \in U_s$, so $(K_0 \setminus \{s\}) \cup \{u\} \in \mathcal{H}_G$. Analogously $(K_0 \setminus \{s\}) \cup \{v\} \in \mathcal{H}_G$. Hence $(K_0 \setminus \{s\}) \cup \{u, v\}$ is a complete subgraph and it can be extended to some clique K . Then $(K_0 \setminus \{s\}) \cup \{u\} \not\subseteq K$ which contradicts the maximality of the clique $(K_0 \setminus \{s\}) \cup \{u\}$. Let now $u \in U_s, v \in U_t$ and $s \neq t$. Then $K_1 = (K_0 \setminus \{s\}) \cup \{u\} \in \mathcal{H}_G, K_2 = (K_0 \setminus \{t\}) \cup \{v\} \in \mathcal{H}_G$. We exchange the cliques K_1, K_2 and the element $t \in K_1 \setminus K_2$. Observe that $K_2 \setminus K_1 = \{v, s\}$. If $K_1 \setminus \{t\} \cup \{s\} \in \mathcal{H}_G$ then $\{s, u\} \in X$ contrary to (10).

Now we shall show a class of graphs in which cliques form e*-semimatroids but not necessarily matroids.

Let n_1, \dots, n_r be a sequence of positive integers such that $n_i \geq n_j$ for $i < j$. Consider a graph $G_{(n_1, \dots, n_r)} = (U_1 \cup U_2 \cup \dots \cup U_r, X)$ where $U_i = \{a_{i1}, \dots, a_{in_i}\}$, $U_i \cap U_j = \emptyset$ for $i \neq j$, $i, j \in \{1, \dots, r\}$. The set X contains all possible edges except those of the following three forms:

$$(11) \quad \{a_{ik_1}, a_{ik_2}\} \text{ for } k_1, k_2 \in \{1, \dots, n_i\},$$

$$(12) \quad \{a_{(2s-1), k'}, a_{2s, k}\} \text{ where } k' \neq k \text{ and } 1 \leq k, k' \leq n_{2s},$$

$$(13) \quad \{a_{(2s-1), t'}, a_{2s, t}\} \text{ where } n_{2s} < t \leq n_{2s-1}, 1 \leq t' \leq n_{2s}, 1 \leq s \leq \lfloor \frac{1}{2}r \rfloor.$$

Observe that the intersection of any clique of the graph $G_{(n_1, \dots, n_r)}$ with any of the sets U_i ($i = 1, \dots, r$) contains exactly one element. Checking that the condition (e^*) holds for the clique is easy.

The following problem is open:

Describe all simple graphs in which cliques form e^ -semimatroids.*

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